

MATH 595 - LECTURE 28

Last time:

- \mathcal{M} = category of smooth manifolds
- **Fibered category** over \mathcal{M} : a functor $\pi: \mathcal{G} \rightarrow \mathcal{M}$ satisfying:
 - (i) For every $f: X' \rightarrow X \notin \mathcal{C}$ in \mathcal{G} over X , pullback exist:

$$\begin{array}{ccc} \exists g & \begin{array}{c} C' \\ \downarrow \\ C \end{array} & \xRightarrow{\pi} \begin{array}{c} X' \\ \downarrow f \\ X = \pi(C) \end{array} \end{array} \quad \text{C}' = f^*C = C|_{X'} \text{ "pullback"}$$

- (ii) Existence of unique lifts:

$$\begin{array}{ccc} \exists' g & \begin{array}{ccc} C_1 & \xrightarrow{\beta_1} & C \\ \downarrow & & \uparrow \\ C_2 & \xrightarrow{\beta_2} & C \end{array} & \xRightarrow{\pi} \begin{array}{ccc} X_1 & \xrightarrow{f_1} & X \\ f \downarrow & & \uparrow \\ X_2 & \xrightarrow{f_2} & X \end{array} \end{array}$$

- **Fiber over X** : $C_X = \{ (C, g) : \pi(C) = X, \pi(g) = id_X \}$ (Gaps)

Fibered categories generalize presheaves

In one direction, given a presheaf (i.e., a contravariant functor):

$$P: \mathcal{M} \rightarrow \text{Sets}$$

one defines a fibered category:

$$\mathcal{G} := \begin{cases} \text{Obj} = \{ (X, \alpha) : X \in \mathcal{M} \neq \alpha \in P(X) \} \\ \text{Arr} = \{ (X, \alpha) \xrightarrow{f} (Y, \beta) \mid f: X \rightarrow Y, P(f)(\beta) = \alpha \} \end{cases}$$

$$\pi: \mathcal{G} \rightarrow \mathcal{M} := \text{forgetful functor}$$

RMK: Note that this is a discretely fibered category, i.e. fibers C_X are identity gaps (\Leftarrow pullbacks are unique)

In the other direction, Given fibred category $\pi: \mathcal{C} \rightarrow \mathcal{M}$ make a choice of pullbacks:

- For each $C \in \mathcal{C}_X$ & map $f: X' \rightarrow X$ choose a lift $\tilde{f}: C' \rightarrow C$

With notation $C' = f^*C$, by property (ii), if $C_1, C_2 \in \mathcal{C}_X$ and $g: C_1 \rightarrow C_2$

Then:

$$\begin{array}{ccc} f^*C_1 & \xrightarrow{f^*g} & f^*C_2 \\ \downarrow & & \downarrow \\ C_1 & \xrightarrow{g} & C_2 \end{array}$$

We obtain a **contravariant map**

$$P: \mathcal{M} \rightarrow \text{Grpds}$$

where:

- $P(X) := C_X$ (A groupoid)
- $P(f: X_1 \rightarrow X_2) := f^*: C_{X_2} \rightarrow C_{X_1}$ (A groupoid morphism)

But this is only a **pseudo functor**: by axiom 2, there is a

$$(f_1 \circ f_2)^* \underset{\tau_{f_1, f_2}}{\cong} f_2^* \circ f_1^*$$

where τ_{f_1, f_2} is a unique natural transformation satisfying some coherence conditions. One calls P a **Lax presheaf of Grpds**. One finds that:

- Different choices of pullbacks \Rightarrow equivalent Lax presheaves of Grpds
- Lax presheaf of Groupoids \Rightarrow Fibred category (similar to above)
- Equivalent Lax presheaves of Grpds \Rightarrow Equivalent Fibred categories (we define equivalence later)

Lax presheaf of groupoids \leftrightarrow Fibred category + choice of pullbacks

Rmk: A choice of pullbacks for $\pi: \mathcal{C} \rightarrow \mathcal{M}$ is also called a **cleavage**.

Examples

1) The fibred category \underline{M} has unique cleavage. The associated presheaf is the strict presheaf:

$$U \mapsto C^\infty(U, M) \quad (\text{identity groupoids})$$

More generally:

strict presheaves of groupoids \leftrightarrow discretely fibred categories
(i.e., unique pullbacks)

2) $\pi: B\mathcal{G} \rightarrow M$ has a natural cleavage: pullbacks of principal \mathcal{G} -bundles

The associated Lax presheaf of groupoids is

$$\cdot X \mapsto \{ \text{Principal } \mathcal{G}\text{-bundles over } M \}$$

$$\cdot (f: X_1 \rightarrow X_2) \mapsto f^* \equiv \text{Pullback of principal } \mathcal{G}\text{-bundles}$$

This is not a strict presheaf:

$$(f_1 \circ f_2)^* P \neq f_2^* (f_1^* P)$$

3) For any Lie group $\mathcal{G} \rightrightarrows X$ we can define a strict presheaf of groupoids:

$$U \mapsto \text{Grpd } \omega \begin{cases} \text{Objs: } C^\infty(U, X) \\ \text{Arr: } C^\infty(U, \mathcal{G}) \end{cases}$$

$$(f: U \rightarrow V) \mapsto \begin{array}{ccc} C^\infty(V, \mathcal{G}) & \xrightarrow{f^*} & C^\infty(U, \mathcal{G}) \\ \downarrow & & \downarrow \\ C^\infty(V, M) & \xrightarrow{f^*} & C^\infty(U, M) \end{array}$$

By correspondence above this defines a discrete fibred category.

When \mathcal{G} is the identity groupoid $X \rightrightarrows X$ we recover \underline{X} .

Note: In general, this fibred category is $\neq B\mathcal{G}$

Prestacks & Stacks: Descent

So far we have not used covering families, i.e., the Grothendieck topology. This allows us to introduce "gluing axioms" similar to sheaves, and define stacks. Recall:

- A covering family of X is $\{U_i \xrightarrow{f_i} X\}$ w/ f_i étale and $\bigcup_i f_i(U_i) = X$.

Consider a presheaf over \mathcal{M} , i.e., a contravariant functor

$$P: \mathcal{M} \rightarrow \text{Sets}$$

Given $f: U \rightarrow V$ & $\alpha \in P(V)$, we use the usual notations,

$$\alpha|_U := P(f)(\alpha)$$

Also, for a covering family $\{f_i: U_i \rightarrow X\}$ we set:

$$U_{ij} := U_i \times_X U_j, \quad U_{ijk} := U_i \times_X U_j \times_X U_k, \text{ etc.}$$

Note that these are all good pullbacks. Recall:

Def:

(i) A presheaf P over \mathcal{M} is separated if

$$\left. \begin{array}{l} \{f_i: U_i \rightarrow X\} \text{ covering family} \\ \alpha, \alpha' \in P(X), \forall i: \alpha|_{U_i} = \alpha'|_{U_i} \end{array} \right\} \Rightarrow \alpha = \alpha'$$

(ii) A presheaf F over \mathcal{M} is a sheaf if

$$\left. \begin{array}{l} \{f_i: U_i \rightarrow X\} \text{ covering family} \\ \alpha_i \in F(U_i), \alpha_i|_{U_{ij}} = \alpha_j|_{U_{ij}} \end{array} \right\} \Rightarrow \exists \alpha \in P(X) \text{ w/ } \alpha|_{U_i} = \alpha_i$$

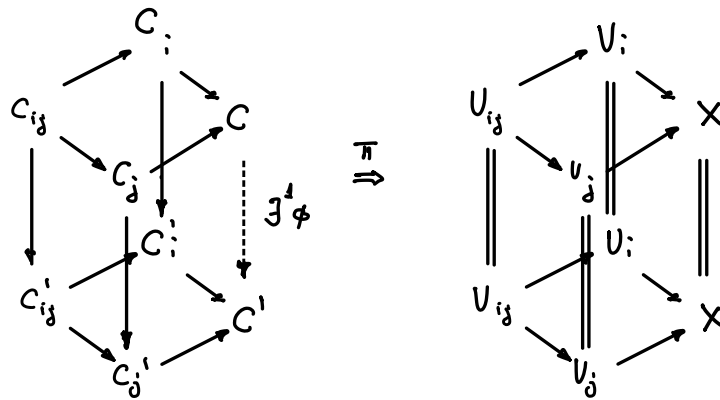
Since fibrewise categories can be thought as lax presheaves one can impose similar properties. These lead to prestacks & stacks.

First we define the analogue of separated:

Def: A fiber category $\pi: \mathcal{C} \rightarrow \mathcal{M}$ is called a prestack if for some cleavage, for any $C, C' \in \mathcal{C}_X$, any covering family $\{f_i: U_i \rightarrow X\}$ and isomorphisms $\phi_i: C|_{U_i} \rightarrow C'|_{U_i}$

$$\phi_i = \phi_j \text{ (on } C|_{U_{ij}}) \Rightarrow \exists^! \text{ isomorphism } \phi: C \rightarrow C'.$$

Bmk: Without a choice of cleavage, the condition is that one can always fill the diagram:



Examples

1) For a Lie group G , $BG \rightarrow \mathcal{M}$ is a prestack; the condition amounts to the gluing axiom for maps of spaces applied to principal bundles.

Similarly, principal G -bundles with connection form a prestack $BG^\nabla \rightarrow \mathcal{M}$: principal bundle connections can also be glued.

More generally, for any Lie groups BG & BG^∇ are prestacks

2) If $\pi: \mathcal{C} \rightarrow \mathcal{M}$ is a discretely fibered category then it is a prestack iff the corresponding presheaf $P: \mathcal{M} \rightarrow \text{Groups}$ is separated.

In particular, for any manifold, \underline{M} is a prestack and for any Lie group G , $X \mapsto C^\infty(X, G)$ gives a prestack.

Def: A prestack $\pi: \mathcal{C} \rightarrow \mathcal{M}$ is called a stack if for some choice of cleavage, for any $X \in \mathcal{M}$, any covering family $\{f_i: U_i \rightarrow X\}$, any $C_i \in \mathcal{C}_{U_i}$ and $\phi_{ji}: C_i|_{U_{ij}} \rightarrow C_j|_{U_{ij}}$, satisfying the cocycle condition

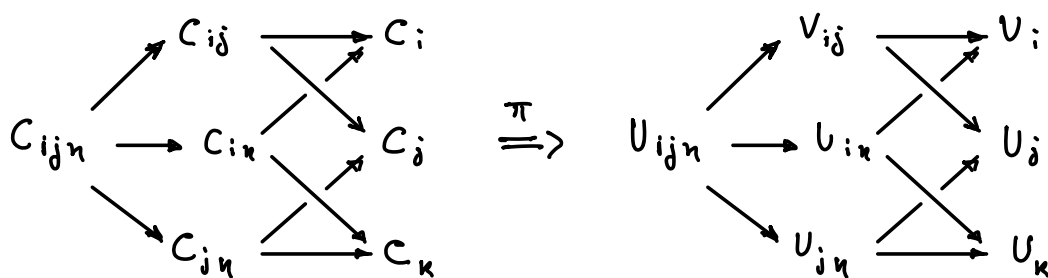
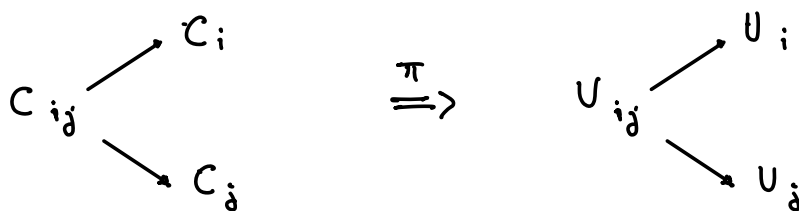
$$\phi_{kj} \circ \phi_{ji} = \phi_{ki} \quad (\text{in } \mathcal{C}_{U_{ijk}})$$

there exists $C \in \mathcal{C}_X$ and isomorphisms $\phi_i: C|_{U_i} \rightarrow C_i$

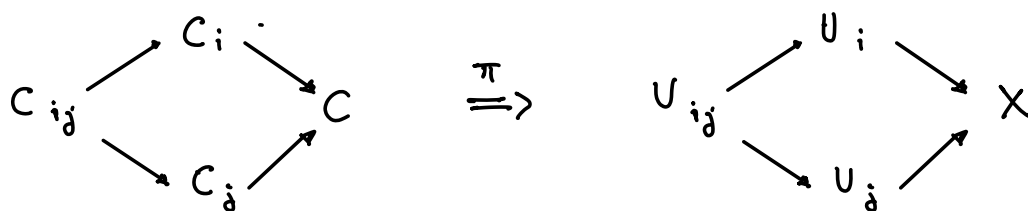
such that

$$\phi_{ji} \circ \phi_i = \phi_j \quad (\text{in } \mathcal{C}|_{U_{ij}})$$

Bmk: Without a choice of cleavage, the condition says that given data:



There is $C \in \mathcal{C}_M$ and arrows $C_i \rightarrow C$ filling diagram:



Bmk: The data $\{C_i, \phi_{ij}\}$ satisfying cocycle condition is call Descent Data. The condition for a prestack to be a stack is called the Descent condition.

Examples:

1) A discretely fibered category is a stack iff the corresponding presheaf is a sheaf. In particular, for any manifold M is a stack.

The fibred category associated w/ $X \mapsto \hat{C}(X, G)$ is not, in general, a stack.

2) For $\pi: BG \rightarrow M$ descent data amounts to a collection of principal G -bundles $P_{U_i} \rightarrow U_i$ and transition functions giving a G -cocycle. Then there exists a principal G -bundle $P \rightarrow M$ and isomorphisms $f_i: P|_{U_i} \cong P_{U_i}$ compatible w/ transition maps. Hence, BG is a stack.

RMK: Given a prestack there is a stackification. If one stackifies the fibred category associated w/ $X \mapsto \hat{C}(X, G)$ one obtains BG .

Maps of Stacks

DEF: Let $\pi_1: C_1 \rightarrow M$ & $\pi_2: C_2 \rightarrow M$ be fibred categories:

i) A map of fibred categories is a functor $\Phi: C_1 \rightarrow C_2$ such that $\pi_2 \circ \Phi = \pi_1$

ii) A 2-isomorphism of fibred categories between $\Phi: C_1 \rightarrow C_2$ and $\Psi: C_1 \rightarrow C_2$ is a natural isomorphism $\tau: \Phi \rightarrow \Psi$ w/ $\tau(x) \in C_x$.

iii) An equivalence of fibred categories is a map $\Phi: C_1 \rightarrow C_2$ admitting a quasi-inverse $\Psi: C_2 \rightarrow C_1$ (so $\Phi \circ \Psi \cong \text{Id}_{C_1}$, $\Psi \circ \Phi \cong \text{Id}_{C_2}$)

A map, 2-isomorphism, or equivalence of stacks is just a map, 2-isomorphism or equivalence of the underlying fibred categories.

Hence, Fibered categories and stacks are both 2-categories. The latter is denoted $\text{St}(\mathcal{M})$.

Properties:

(i) If $\underline{\Phi}: \mathcal{C}_1 \rightarrow \mathcal{C}_2$ is a map of fibered categories then $\underline{\Phi}$ is an equivalence iff for every $X \in \mathcal{M}$ the restriction to the fibers $\underline{\Phi}|_X: \mathcal{C}_1|_X \rightarrow \mathcal{C}_2|_X$ is an equivalence of categories.

(ii) If $\underline{\Phi}: \mathcal{C}_1 \rightarrow \mathcal{C}_2$ is an equivalence of fibered categories then \mathcal{C}_1 is a prestack (resp. stack) iff \mathcal{C}_2 is a prestack (resp. stack).

Examples

1) Given $\phi: M \rightarrow N$ we obtain a map of stacks:

$$\underline{\phi}: \underline{M} \rightarrow \underline{N} \quad \left\{ \begin{array}{l} (f: X \rightarrow M) \mapsto \phi \circ f: X \rightarrow N \\ \begin{array}{ccc} X_1 & \xrightarrow{f_1} & \\ \downarrow g & \searrow & \\ X_2 & \xrightarrow{f_2} & \end{array} \mapsto \begin{array}{ccc} X_1 & \xrightarrow{\phi \circ f_1} & N \\ \downarrow g & \searrow & \\ X_2 & \xrightarrow{\phi \circ f_2} & \end{array} \end{array} \right.$$

Every map of stacks $\underline{\Phi}: \underline{M} \rightarrow \underline{N}$ is of this form. Note that

$$\underline{M} \text{ is discretely fibered} \Rightarrow \begin{cases} \text{No non-trivial 2-isomorphisms} \\ \text{between maps} \end{cases}$$

So we have a full embedding of 2-categories:

$$\mathcal{M} \rightarrow \text{St}(\mathcal{M}) \quad \left\{ \begin{array}{l} M \mapsto \underline{M} \\ \phi \mapsto \underline{\phi} \end{array} \right.$$

Remark: One can also show that for any stack $\pi: G \rightarrow \mathcal{M}$ there is a canonical equivalence of groupoids:

$$G_X \simeq \text{Hom}(X, G)$$

This example and remark formalize the idea that we can understand a (generalized) space by looking at all maps into the space.

2) Let $\Phi: \mathcal{G} \rightarrow \mathcal{H}$ be a morphism of Lie groupoids. One obtains a map of stacks $\Phi_*: \mathcal{B}\mathcal{G} \rightarrow \mathcal{B}\mathcal{H}$ as follows:

$$\cdot \quad \begin{array}{ccc} \mathcal{G} & \downarrow & \mathcal{P} \\ \mathcal{G} & \rightrightarrows & \mathcal{M} \end{array} \quad \longmapsto \quad \begin{array}{ccc} \mathcal{G} & \downarrow & \Phi_*(\mathcal{P}) = (\mathcal{H} \times_N \mathcal{P}) / \mathcal{G} \\ \mathcal{H} & \rightrightarrows & \mathcal{N} \end{array}$$

$$\cdot \quad (f: \mathcal{P}_1 \rightarrow \mathcal{P}_2) \longmapsto \Phi_*(f)([h, p]) = [h, f(p)]$$

Exercise: Show that $\Phi_*: \mathcal{B}\mathcal{G} \rightarrow \mathcal{B}\mathcal{H}$ is an equivalence of stacks iff $\Phi: \mathcal{G} \rightarrow \mathcal{H}$ is a Morita map.

Def. Let $\Phi: \mathcal{G}_1 \rightarrow \mathcal{G}_2$ be a map of fibered categories over \mathcal{M} .

(i) Φ is a monomorphism if for every $X \in \mathcal{M}$ the restriction to the fibers $\Phi: (\mathcal{G}_1)_X \rightarrow (\mathcal{G}_2)_X$ is fully faithful.

(ii) Φ is a epimorphism if for every $X \in \mathcal{M}$ and $C_2 \in (\mathcal{G}_2)_X$ there is a covering family $\{f_i: U_i \rightarrow X\}$ and $C_i \in (\mathcal{G}_1)_{U_i}$ such that $\Phi(C_i) \simeq C_2|_{U_i}$.

Rmk: If $\Phi: (\mathcal{G}_1)_X \rightarrow (\mathcal{G}_2)_X$ is ess. surjective for every $X \in \mathcal{M}$, then (ii) holds for trivial covering family $\{\text{id}: X \rightarrow X\}$. So (ii) is weaker than this condition. Hence, an equivalence of fibered categories is both a monomorphism and an epimorphism but not the converse.

Exercise: Let $\phi: M \rightarrow N$ be a smooth map. Show that

for the map of stacks $\underline{\phi}: \underline{M} \rightarrow \underline{N}$:

(i) $\underline{\phi}$ is always a monomorphism;

(ii) $\underline{\phi}$ is epimorphism iff ϕ is a surjective submersion.