MATH 595 - LECTURE 23

Last time : DIFFERENT Noticous of Equivalences

- A Lie Gnocpois Konphism & D S SP is called A: H H H H + N
- · looncaphian : J E: & g a/ E. E. idg, Do E = idg
- · Strows Equivaliance: 三臣: 部一日山山 里·西 ? idg, 豆 平 idg
- · Monita HAP :

Wo saw That:

Isonoophism => strong Equivalence => Monita map Monita is The Right Notion of Equivalence:

Theoren (del Hogo)

A GAOUPOID HOAPHIGH E: G - H is a Monita Map ipp it presences transvensal Bata

PROOF :

(1)
$$\overline{\Phi}$$
 Fully Faith Full =>

$$\begin{cases}
G_{n} & fl_{\varphi(n)} \\
M/g & N/g \text{ injective} \\
d_{1}\varphi : \mathcal{V}_{x}(G) & \mathcal{V}_{\varphi(i)}(\varphi(G)) \text{ injective}
\end{cases}$$

Since I is cateconical equivalence (set Theoretically) Frast two follows

For last one, use The pullback Dia CRAM:

Hince :

On the other haws, The Diacram:

$$T (\mathcal{H} \times M) \xrightarrow{dpe_i} T_h \mathcal{H} \xrightarrow{dt} T_N$$

$$\stackrel{(h,w)}{\longrightarrow} dpe_i \downarrow \qquad ds \downarrow \qquad \stackrel{tin}{\longrightarrow} \forall d\lambda_{h'}$$

$$T_M \xrightarrow{dq} T_N \longrightarrow \forall (q(0))$$

Shows that

$$d_{\alpha}(tope_{\star}) \quad \text{subjective} \quad <=> \quad d\phi: T_{\alpha} M \rightarrow V_{\phi(s)}(\phi(G)) \quad \text{subjective}$$

$$(=> \quad cl\phi: V_{\alpha}(G) \rightarrow V_{\phi(s)}(\phi(G)) \quad \text{subjective}$$

(III)] prescrives transucesal sata =>] Folly FAITHFUL

$$G_m \simeq \partial d_{q(n)}$$

 $M/g \sim N/gg$
 $M/g \sim N/gg$
 $G_n \simeq \partial d_{q(n)}$
 $G_n \simeq \partial d_{q(n)}$

$$d_{*}\varphi: U_{n}(\mathcal{O}) \rightarrow U_{\varphi_{l}u_{l}}(\varphi(\mathcal{O}) \quad \text{sugardius} \Rightarrow \begin{cases} \varphi \times \varphi: M \times M \rightarrow N \times N \\ t \times s: S \end{pmatrix} \rightarrow N \times N \\ \text{haps} \end{cases}$$

$$= \sum_{k=1}^{k} \begin{cases} u_{k}(\mathcal{O}) & u_{k}(\mathcal{O}) \\ u_{k}(\mathcal{O}) & u_{k}(\mathcal{O}) \\ u_{k}(\mathcal{O}) & u_{k}(\mathcal{O}) \end{cases}$$

txs $M \times M \xrightarrow{4 \times 0} N \times N$ H REMAINS to check that diff. or $G \xrightarrow{-} N \times M \times BR$ is bijective NXN

This Follows FROM DINGRAM CHABING:

$$0 \longrightarrow \mathcal{T}_{g}(\overline{t}^{(y)} \cap \overline{s}^{(w)}) \longrightarrow \mathcal{T}_{g}\zeta_{g} \xrightarrow{dt \times ds} \mathcal{T}_{w} M \times \mathcal{T}_{y} M \longrightarrow \mathcal{V}_{g}(G) \longrightarrow O$$

$$d_{g} \overline{\Phi} \downarrow s_{1} \qquad d_{g} \overline{\Phi} \downarrow \qquad \varphi \times \varphi \downarrow \qquad s_{1} \downarrow d\varphi$$

$$0 \longrightarrow \mathcal{T}_{g}(\overline{t}^{(y)} (\varphi(s_{1})) \cap \overline{s}^{(y)}) \longrightarrow \mathcal{T}_{g}\zeta_{g} \xrightarrow{dt \times ds} \mathcal{T}_{N} \times \mathcal{T}_{N} \longrightarrow \mathcal{V}_{g}(\varphi(b)) \longrightarrow O$$

Examples

1) IF G = M is a transitive Lie croopers, then For any pGM, $G_P \subseteq G$ is a Monita Map: $M/g = 4 \times 3$ # $U_p(M) = 0$

2) IF G = M is my Groepoid and TCM is a submanifold That intenseels thransversally every orbit or G, then $G|_{=}=T$ is a Lib Choupeid and $G|_{=}=G$ is a Monita Hap.:

• (Π, \mathcal{F}) Foliation: $Hol(\Pi, \mathcal{F})|_{T} \Rightarrow T \subset Hol(\Pi, \mathcal{F}) \Rightarrow M$ is Honita • $G \ltimes \Pi \Rightarrow \Pi$ action: $(G \ltimes \Pi)|_{T} \Rightarrow T \subset (G \ltimes \Pi) \Rightarrow \Pi$ is Florita.

3) Composition of Monita maps is a Monita Map

<u>Def</u>: $G_1 = \Pi_1 \notin G_2 = \Pi_2$ and Herita equivalent if $\exists S = M$ # Monita maps $\overline{\sigma}_1 : \overline{\sigma}_1$



To prove This is An Equivalore Colation, as NGED WEAK pellbach or grocpoiss:

Kincas:

- Objects or
$$\mathcal{F} \times \mathcal{F} := Q \times \mathcal{G} \times \mathcal{R}$$
, i.e.:

$$\begin{array}{c} g\\ (q, g, r) & \text{with} & f \times (r) = S(g) \\ g(q) = t(g) \end{array}$$

- ARROWS OF PX X = ARROW FROM (9,,91,V1) to (9,,92,V2) is A pain (h, K) such that

Hence, we can also identify arrows $\omega/(h, g, K) \in \mathcal{H} \times \mathcal{G} \times \mathcal{H}$ with: $\begin{cases} s(g) = \psi(t(K)) \\ t(g) = \varphi(s(h)) \end{cases}$

The Diagram (*) is not connectative: (***) sage that $id: G \rightarrow G$ is a Natural isomorphism between two scores of (*). In other works, The Diagram (*) is weakly connectative. IN GENERAL, WEAK Pullback is not snooth. Proposition

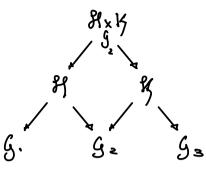
IF The Fiberco product $Q \times G \times R$ exists, then the weak pullback is a Lie Chooperd. This is the case, e.g., if $\overline{\Phi}$ (resp., $\overline{\Psi}$) is A Monita map, and in this case the base chance of $\overline{\Psi}$ (resp. $\overline{\Psi}$) is a florita map

Proof: The First part Follows From Diagram:

For The secon part, observe That is I to Tonita, then

SO WEAK pullbrok is swooth by First part. Exencise: Verify That à is Monita.

<u>Concllary</u>: Monita Equivalence is an equivalence eclation <u>Proof</u>: Reflexive and symmetry is obvious. Fa transitive, by Proposition, we can use weak pullback.



AND THE FACT That composition of Nonita raps is Monita.

Examples.

1) A Lie cacupero G = M is Monita equivalent to A Manifelo N = N iff Z surgestive submension $\phi: M \rightarrow N \notin A$ Groupero isomorphism $G \simeq M \times M$.

2) An Action CAOLDERS GRM = M is Monita équivalent to A MANIFELO N=N IFF The Action is proper & FREF.

3) A lio gnorpeis G = M is Monita Equivalent to A lio gnorp G iFF et is transitive and any of its isoloopy Gnoops is isomorphic to G.

4) A Lie gnoupers G = M is Monita equivalent to a Discrete Gnoup iff it is transitive of any up its restropy choops is Discrete IFF txs: $G - M \times \Pi$ is a covering.

We will discuss later abat properties of a lie Croupers Are invariant unser Monita Equivalence.

Monita Equivalence can also be seen as isonorphisms n A contain category: Der. A Generalizes Monphism E/ : G -- + H is given by a pair of Lie croepers reaphisms:

Fro generalizes Monphisms <u>I, I</u>: G H ADE iDentical

if There is a Third Generalized Konphism 43/2: G---> fl Fitting into a Weakly connetative Deagram:

(This is an equivalence relation on the set of GEN. nonphisms)

Monita equivalences = generalizes iscrenphisms (inventible Generalizes) Rorphisms

Exercise: Show Mat

(i) eveny concratizes map $\frac{\underline{w}}{\underline{w}}$: $G_1 \cdots = G_2$ has a Ropresculative $\frac{\underline{w}}{\underline{w}}$: $G_1 \cdots = G_2$, where $\underline{\Phi}'$ covers a surgestive submension

(ii) Eucry Monita Equivalence can be represented by a pair of Monita maps which are surjective submensions.

<u>Hint</u>: Given Monita map $\overline{\Phi}: \mathcal{H} \to \mathcal{G}$ consider weak pullback $\mathcal{G}_{X}\mathcal{H} \to \mathcal{H}$