MATH 595 - LECTURE 11

LAST TIME : GIUCN Lie AlGEBADIS A-M

- $\frac{A c_{\text{CNNGetion on U.b.}}{E \to M}}{\nabla_{f_x} S = f \nabla_{f_y} S} = \int \nabla_{f_y} S + \int \nabla_$
- A = path: $a: I \rightarrow A \quad s.t. \quad p(a(t)) = \frac{d}{dt} \mathcal{F}_{a}(t) \quad (\mathcal{F}_{a} = base path of a)$

<=> adt : TI - A Lie alsebrois nonphism

· <u>A-DENIVATIOS</u>: A-path a: I→E path u: I→E above &a: I→M J Dau: I→E path above &a

Geonetrice interpotation of Curvature

With base map &: IXI -> M (paranctenized surface of A-paths) Fix path u: IXI -> E Above &.

$$\begin{array}{l} \rho \circ \overline{\Phi} = d & \Rightarrow \\ \left\{ \begin{array}{l} \rho \circ \overline{\Phi}_{1}\left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial t}\left(t,\varepsilon\right) \\ \rho \circ \overline{\Phi}_{2}\left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \end{array} \right. \\ \left. \left. \begin{array}{l} \rho \circ \overline{\Phi}_{2}\left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \\ \rho \circ \overline{\Phi}_{2}\left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \end{array} \right. \\ \left. \left. \begin{array}{l} \varphi \circ \overline{\Phi}_{2}\left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \\ \left. \left. \begin{array}{l} \varphi \circ \overline{\Phi}_{2}\left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \\ \left. \left. \begin{array}{l} \varphi \circ \overline{\Phi}_{2}\left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \\ \left. \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \right) \left(t,\varepsilon\right) \\ \left. \left(t,\varepsilon\right) = \frac{\partial \overline{e}}{\partial \varepsilon}\left(t,\varepsilon\right) \\ \left(t,$$

~ A-Derivatice $D_{\underline{\mathfrak{F}}_2}$ h: $[0,1]\times[0,1] \rightarrow E$

Proposition With These Notations:

$$\mathcal{R}^{\nabla}(\Phi_1, \Phi_2) \mathcal{U} = \mathcal{D}_{\Phi_1} \mathcal{D}_{\Phi_2} \mathcal{U} - \mathcal{D}_{\Phi_2} \mathcal{D}_{\Phi_1} \mathcal{U}$$

We can pullback the connection via the Algebrois MAP = T(IXI) - A, AND NO FIND:

$$D_{\overline{\Phi}_1} u = \nabla_{\underline{d}_1} u \qquad D_{\overline{\Phi}_2} u = \nabla_{\underline{d}_2} u \qquad d_{\overline{d}_2} u$$

The coevalues are related by pullbach: $R^{\nabla'}(\underline{\Phi}^{*}\alpha, \underline{\Phi}^{*}\beta)(\underline{\mathcal{X}}^{*}S) = \underline{\mathcal{X}}^{*}(R^{\nabla}(\alpha, \beta)(S))$ $\Rightarrow R^{\nabla}(\underline{\Phi}_{1}, \underline{\Phi}_{2})M = R^{\nabla'}(\underline{d}_{1}, \underline{d}_{2})M$ Since $[\underline{d}_{4t}, \underline{d}_{2}]zo$, the result fellows.

Georginic Interpotation of Tonsion

Consider a vector bundle map

$$T(\mathbf{I}\times\mathbf{I}) \xrightarrow{\Phi} A \qquad \Phi = \overline{\Phi}, dt + \overline{\Phi}, dt$$

$$\downarrow \qquad \downarrow \qquad \downarrow \qquad \downarrow$$

$$\mathbf{I}\times\mathbf{I} \xrightarrow{\bullet} M \qquad \overleftarrow{\bullet} \qquad \overleftarrow$$

Assume that I preserves anchors:

$$6 \circ \overline{\Phi} = q$$

tale still have the A-paths obtained by PREEZING E/t AS BEFAG. (proservation of brackets was not used)

Proposition

Given AN A-connection V on A, A V.b. MAP \$:T(IXI) -A compatible top Anohons is a Lie algebroid reaphiem iff

$$T^{\nabla}(\Phi^{1},\Phi^{2}) = D^{\Phi}\Phi^{2} - D^{\Phi}\Phi^{1}$$

PADOF.

(=>) Pullbach ∇ alove $\underline{\Phi}$ to obtain classical connection an T(IXI)And apply classical Result.

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(=) We will concease to This reacht and its proof later.

PARALLEL TRANSport

<u>DEF.</u> Given $E \rightarrow th$ with A-connection ∇ we say that $u: I \rightarrow E$ is <u>parallel</u> along an A-path $a: I \rightarrow A$ if a covere \mathcal{J}_A and: $D_a u = 0$

As in classical case, one can 11-transport rectors:

<u>Prop.</u> Given (E, ∇) and A-path $a: I \rightarrow A$, For each $u_0 \in E_{talos}$ There exists a unique // cueve $u: I \rightarrow E$ along a and $u(o)=u_0$. Moneover, the cub-point $u(\Delta)$ depends linearly on u_0 . Proof.

Pullback ∇ plong A-path $a: I \rightarrow A$ to obtain classical connection on $\mathcal{S}_{a}^{*} E \rightarrow [0,1]$. This reputes result to existing of //-thouspost for classical connection on [0,1].

Altermative: In local coencirvates, equation For /1-transport is:

$$\begin{cases}
\frac{d}{dt} u^{k}(t) = -\int_{es}^{t} (v_{a}(t)) u^{t}(t) u^{s}(t) \\
u(0) = u_{0}
\end{cases}$$

linear t-offendent ODS => solutions exist Forall t

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11-transport map along A-path:

Examples

$$T_{a}: V(G) \rightarrow V_{\delta_{a}(i)}(G)$$
 liven A-holonomy of O

3.
$$A = \underline{a} \rightarrow 4x \underline{b}$$
, $V \in \operatorname{Rep}(\underline{a})$
 $A : \underline{I} \rightarrow \underline{a}$ any path $\Longrightarrow \widetilde{c}_{a} : V \rightarrow V$
What is This Amap?
 $dd, \dots, d_{d} \underline{b}$ basis for \underline{a} ; $de_{\underline{a}}, \dots, e_{\underline{k}} \underline{b}$ basis For V
 $e(d_{1})(e_{r}) = A_{ir}^{s}e_{\underline{a}}$ $a(t) = a^{i}(t)a_{i}$ $u(t) = u^{s}(t)e_{\underline{s}}$
 $(D_{\underline{a}} u)(t) = (\frac{du^{s}}{dt}(t) + A_{ir}^{s} n^{i}(t)u^{s}(t))e_{\underline{s}}$
So $T_{\underline{a}}(v) = u(t)$ above :
 $\int \frac{du}{dt}(t) = -e(a(t))u(t)$

Internating to gnoup eep $G \rightarrow G(V)$ we obtain: $\hat{U}_{a}(v) = \exp(-\int_{a}^{1} (t) dt) \cdot v$

GROWMANY /1-transport satisfies 3 basic properties: (i) Invariance under Reparameterization or paths (ii) Concatenation of paths = composition: $\Sigma_{y_1,y_2} = \Sigma_{y_1} \cdot \Sigma_{y_2}$ (iii) Invariance under path-homotopy when connection is plat

We Extand ALL these properties For A - connections.

<u>Paop</u>. Let $\alpha: I \rightarrow A$ be A-path and $\overline{U}: [0,1] \rightarrow [0,1], \overline{U}(1=0,\overline{U}(1))$. Set $\alpha^{\overline{U}}(1):= \overline{U}(1) \alpha(\overline{U}(1))$ Then $\alpha^{\overline{U}}: [0,1] \rightarrow A$ is A-path and For any A-connection //-1 hansport along $\alpha = //-1$ hansport along $\alpha^{\overline{U}}$ $\frac{P_{noof.}}{P(a^{T}(t))} = \tilde{T}(t) P(a(T(t))) = \tilde{T}(t) P_{a}(T(t)) = \frac{d}{dt} V_{a}(T(t))$ $P(a^{T}(t)) = \tilde{T}(t) P(a(T(t))) = \tilde{T}(t) P_{a}(T(t)) = \frac{d}{dt} V_{a}(T(t))$ $So a^{T}(t) \text{ is } A \text{-path}.$

To perform //-transport along a we can pull-back along a: TI - A and Deduce to enoinary //-transport For pullback dunale J'E - I along id-path M(t)=t

For reparametenizes A-path, we have similar Provedice. But Then two reductions are just 11-transport For TI-econoscilien Pelative to A and Reparametenization MOT. By codulary result, The 11-transport Maps coincide.

• Using Reparameterisations that satisfy $T^{(n)}(0) = T^{(n)}(1) = 0$, $\forall m \in \mathbb{N}$ we obtain A-path $A^{\tilde{L}}$ with $A(0) = O_m A(1) = O_y$ and all its Dividential at two, 1 vanish.

• $A_1 : I \rightarrow A \notin A_2 : I \rightarrow A \iff A | A | Divided integration of teo, 1$ If $\partial_{A_1}(0) = \partial_{A_2}(0)$: $A_1 \circ A_2(t) = \begin{cases} 2A_2(2t), 0 \le t \le 1/2 \\ 2A_1(2t-1), 1/2 \le t \le 1 \end{cases}$

<u>Proposition</u>. UNDER These conditions, For Any A-econnection: $T_{a_1 \circ a_2} = T_{a_1} \circ T_{a_2}$

PROOF. Exencise.

Der. Two A-paths Q., Q: I - A ANE CALLED <u>A-path honotopic</u> if There exists A Lie Algebroid Morphism

$$\Phi: T(I \times I) \rightarrow A, \quad \Phi = \Phi_1 dt + \Phi_2 dt$$

$$\Phi_1(t, 0) = A_0(t), \quad \Phi_1(t, 1) = A_1(t)$$

$$\Phi_2(0, t) = \Phi_2(0, t) = 0$$

RmK.

・ Since 重 is Alecensia monphism, it satisfies: Po更=dを

where & : [0,1] × [0,1] - M is base path or \$. As we saw before :

For each
$$\mathcal{E}$$
: $t \mapsto \overline{\Phi}_{1}(t, \varepsilon)$ is A-path
For each t : $\mathcal{E} \mapsto \overline{\Phi}_{2}(t, \varepsilon)$ is A-path

$$\begin{cases} \overline{\Phi}_{2}(0,e)=0 \implies \partial \mathcal{F}(0,e)=0 \implies \mathcal{F}(0,e) \text{ constant} \\ \overline{\partial e} \end{cases}$$

$$\begin{cases} \overline{\Phi}_{2}(1,e)=0 \implies \partial \mathcal{F}(1,e)=0 \implies \mathcal{F}(1,e) \text{ constant} \\ \overline{\partial e} \end{cases}$$

So Qo & Q, have path-honotopic base paths

<u>Proposition</u> For a Flat A-connection, Given any two A-path horotopic pathe a. & a. ;

$$l_{a_o} = l_{a_o}$$

$$\frac{\underline{P}nooF}{-\overline{\Phi}: T(I \times I) \rightarrow A \text{ bs an } A \text{ path honolopy bolderu } Q_{e} \neq Q_{e}, \\ - \widehat{C}_{\overline{\Phi}_{i}}^{t,\epsilon}: E_{\sigma(0,0)} \rightarrow E_{\sigma(t,\epsilon)} //-thenespect \text{ alones } t \mapsto \overline{\Phi}_{i}(t,\epsilon) \\ E_{\sigma(0,\epsilon)}^{n} \\ Fon my \quad u \in E_{\sigma(0,0)}, \quad C(t,\epsilon) := T_{\overline{\Phi}_{i}}^{t,\epsilon}(u) \quad gives \text{ a } Cuous \\ above \quad \forall : I \times I \rightarrow \Pi, \quad \text{satisfyind}: \quad D_{C} = O \\ \overline{\Phi}_{i} \\ Claim. \\ D_{\overline{\Phi}_{2}}^{c} = O \\ Assum(NS \quad Claim, \ \omega e \quad have: \\ \frac{d}{d\epsilon} e(a,\epsilon) = (D_{\underline{\Phi}_{2}}^{c})(a,\epsilon) = O \\ \end{array}$$

So $C(1, e) = T_{\overline{\Phi}_{1}}^{1, e}(n)$ is constant, showing that $T_{a_{0}}(u) = T_{\overline{\Phi}_{1}}^{1, 0}(u) = T_{\overline{\Phi}_{1}}^{1, 1}(u) = T_{a_{1}}(u)$

PROOF OF CLAIM:

At t=0, boundary conditions for A-path honotopy $\overline{\Phi}$ give: D=c (0,e) = d c(0,e) = 0

$$D_{\Phi_2} c (o, e) = \frac{d}{de} c (o, e) = c$$

Since V is Flat, Georetaic interpretation of curvature Gives.

$$D_{\underline{\Phi}_1} D_{\underline{\Phi}_2} c = D_{\underline{\Phi}_2} D_{\underline{\Phi}_2} c = 0$$

So $D_{\Phi_2}C$ is parallel along $t \mapsto \Phi_1(t, \varepsilon)$. By uniqueness of 11-paths, we must have $D_{\Phi_2}C = 0$