# Principles of Safe Autonomy ECE 484 Lecture 2: System Safety 

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## Welcome from Safe Autonomy team!

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## Last time on how to assure safety of an autonomous system

## "Testing can be used to show the presence of bugs, but never to show their absence!" --- Edsger W. Dijkstra

Because there are infinitely many executions and we can only test finitely many of those in any testing algorithm

In a probabilistic sense also, purely using data to gain safety assurance is not practical
Data required to guarantee a probability of $10^{-9}$ fatality per hour of driving is proportional to its inverse, $10^{9}$ hours, 30 billion miles

To learn or extrapolate about all---infinitely many---executions from a finite sampling of executions, we need to make some assumptions about the system. A collection of these assumptions defines a model

On a Formal Model of Safe and Scalable Self-driving Cars by
Shai Shalev-Shwartz, Shaked Shammah, Amnon Shashua, 2017
(Responsibility Sensitive Safety)

## Roadmap

- A simple class of models: automata.
- What are executions of automata: sequence of states
- What are requirements?
- Reachable states, why we care to compute and why that can be hard
-Invariants as approximations of reachable states

Model (switch to notes)


## A "simple" safety scenario

A car moving down a straight road has to detect any pedestrian in front of it and stop before it collides.


Not a trivial requirement


## "simple" =三 Easy



## MPO: Simulate model for testing


"All models are wrong, some are useful."

Wrong and useless models


FIGURE 4. A turkey using "evidence"; unaware of Thanksgiving, it is making "rigorous" future projections based on the past. Credit: George Nasr

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HIGHI.Y IMPHOBABLE

Nassim Nicholas Taleb

## Baked-in Assumptions in our example

- Perception.
- Sensor detects obstacle iff distance $d \leq D_{\text {sense }}$


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- No false positives, negatives, probabilities
- Pedestrian is known to be moving with constant velocity from initial position. This will be used in the safety analysis, but not in the vehicle's automatic braking algorithm
- No sensing-computation-actuation delay.
- The time step in which $d \leq D_{\text {sense }}$ becomes smaller is exactly when the velocity starts to decrease



## Baked-in Assumptions (continued)

- Mechanical or Dynamical assumptions
- Vehicle and pedestrian moving in 1-D lane.

- Does not go backwards.
- Perfect discrete kinematic model for velocity and acceleration.
- Nature of time
- Discrete steps. Each execution of the above function models advancement of time by 1 step. If 1 step $=1$ second, $x_{1}(t+1)=$ $x_{1}(t)+v_{1}(t) .1$
- We cannot talk about what happens between [t, t+1]
- Atomic steps. 1 step = complete (atomic) execution of the program.
- We cannot directly talk about the states visited after partial execution of program


## Summary

- Absolute safety checking boils down to showing that none of the executions of the automaton reaches an unsafe set $U$
- To reason about all executions of we have to work with infinite sets of states
- One way to compute infinite sets is using the Post operator
- But, computing all executions for unbounded time can be hard
- If we can guess an invariant satisfying conditions of Proposition 1.1, that can give a shortcut for proving safety
- The inavariant may contain important information about conserved quantities, and thus, may tell us why the system is safe, and not just that it is so
- Mind the gap between model and reality
- Next. Application of invariants in braking example

