

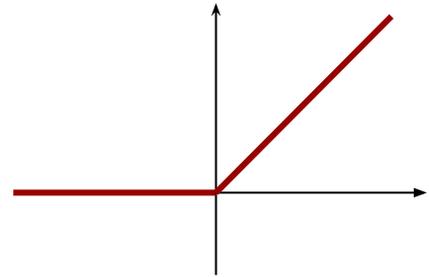
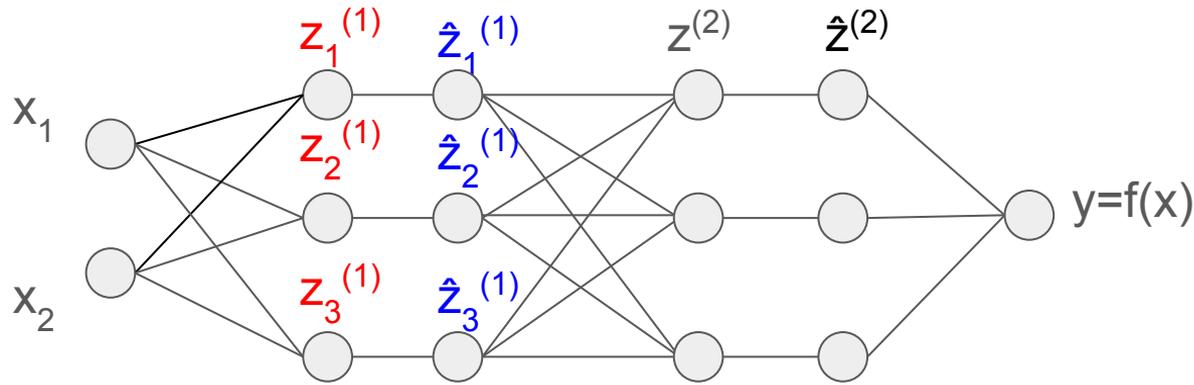
ECE598HZ: Advanced Topics in Machine Learning and Formal Methods

# Lecture 10: Neural Network Verification with Bound Propagation Algorithms (Part I)

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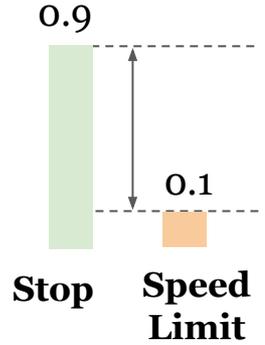
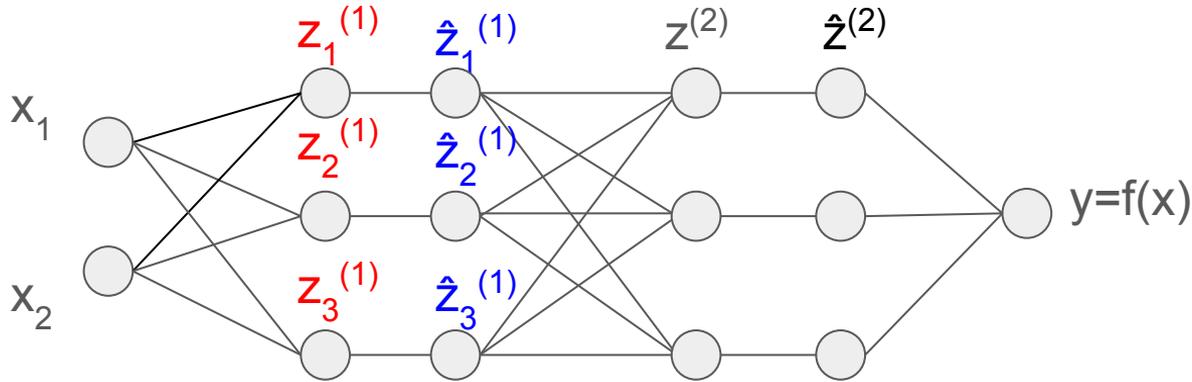
# Review: Neural Networks (NNs)



Linear layers:  $z^{(1)} = W^{(1)} x$        $z^{(2)} = W^{(2)} \hat{z}^{(1)}$        $y = w^{(3)T} \hat{z}^{(2)}$

Nonlinear layers:  $\hat{z}_j^{(i)} = \sigma(z_j^{(i)})$  (assume  $\sigma$  is ReLU for now)

# Review: NN verification as an **optimization** problem



Does there  $\exists x$ , s.t.  $x \in S \wedge y \leq 0 \wedge y = f(x)$

Input domain under consideration

Negation of the desired property

$$y^* = \min_{x \in S} f(x)$$

MILP and LP

# Review: stable vs. unstable neurons

$$\hat{z}_j^{(i)} \leq z_j^{(i)} - l_j^{(i)}(1 - p_j^{(i)})$$

$$\hat{z}_j^{(i)} \leq u_j^{(i)} p_j^{(i)}$$

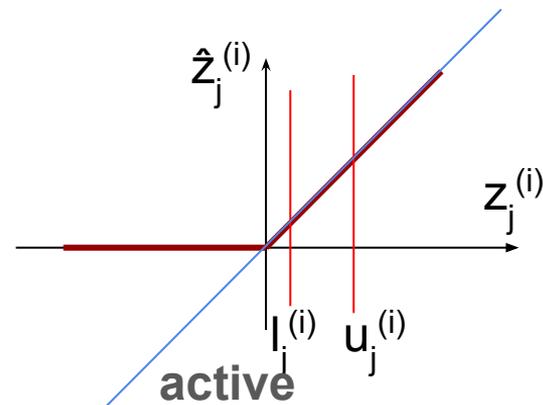
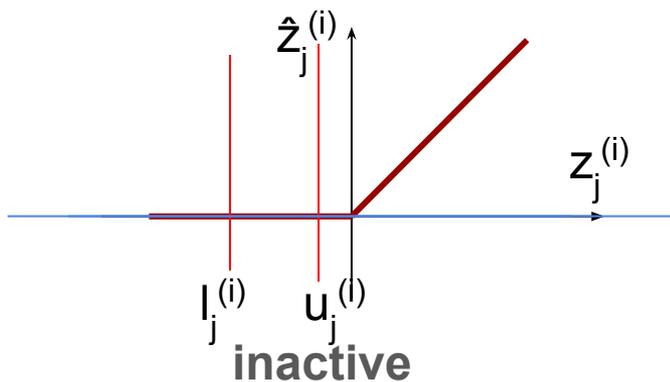
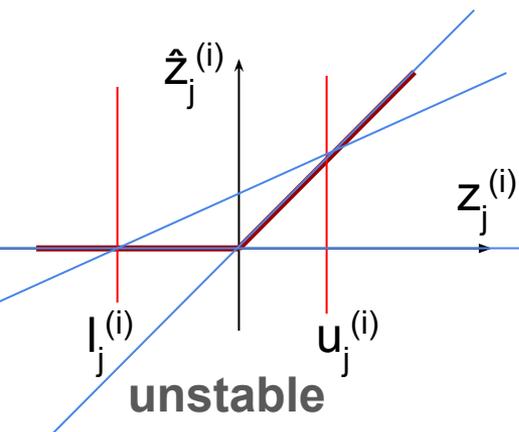
$$\hat{z}_j^{(i)} \geq z_j^{(i)}$$

$$\hat{z}_j^{(i)} \geq 0$$

$$0 \leq p_j^{(i)} \leq 1$$

$$\hat{z}_j^{(i)} = 0$$

$$\hat{z}_j^{(i)} = z_j^{(i)}$$



# Review: triangle relaxation for unstable ReLU neurons

Each ReLU is represented by

$$\hat{z}_j^{(i)} \leq z_j^{(i)} - l_j^{(i)} (1 - p_j^{(i)})$$

$$\hat{z}_j^{(i)} \leq u_j^{(i)} p_j^{(i)}$$

$$\hat{z}_j^{(i)} \geq z_j^{(i)}$$

$$\hat{z}_j^{(i)} \geq 0$$

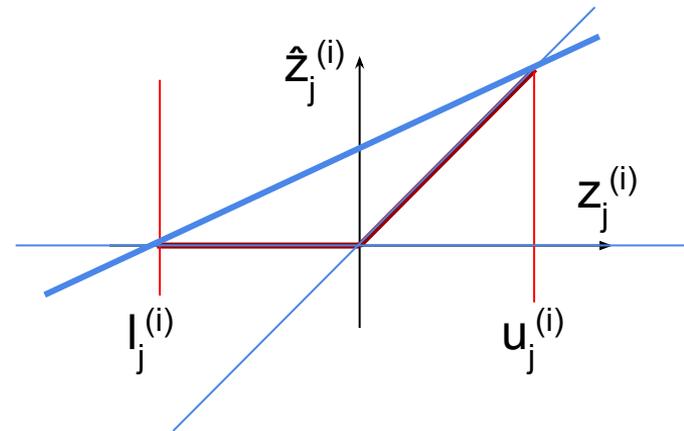


$$\hat{z}_j^{(i)} \leq \frac{u_j^{(i)}}{u_j^{(i)} - l_j^{(i)}} z_j^{(i)} - \frac{u_j^{(i)} l_j^{(i)}}{u_j^{(i)} - l_j^{(i)}}$$

$$\hat{z}_j^{(i)} \geq z_j^{(i)}$$

$$\hat{z}_j^{(i)} \geq 0$$

“Triangle” relaxation



# Today: more efficient algorithms for NN verification

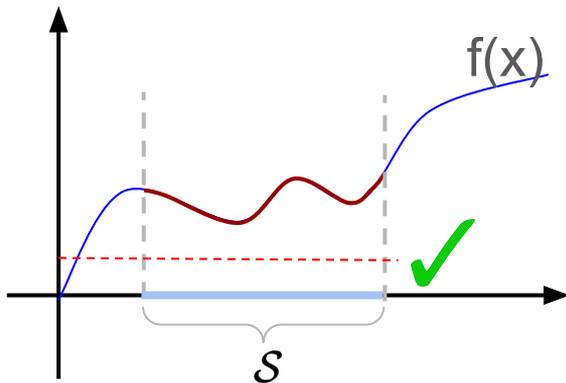
Solving neural network verification using SMT solvers

Solving neural network verification using optimization (MIP/LP)

Solving neural network verification using **bound propagation (this lecture!)**

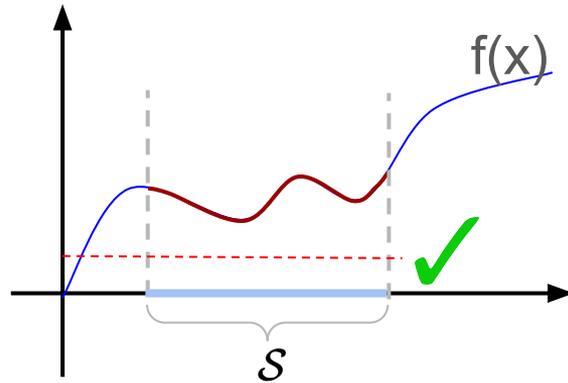
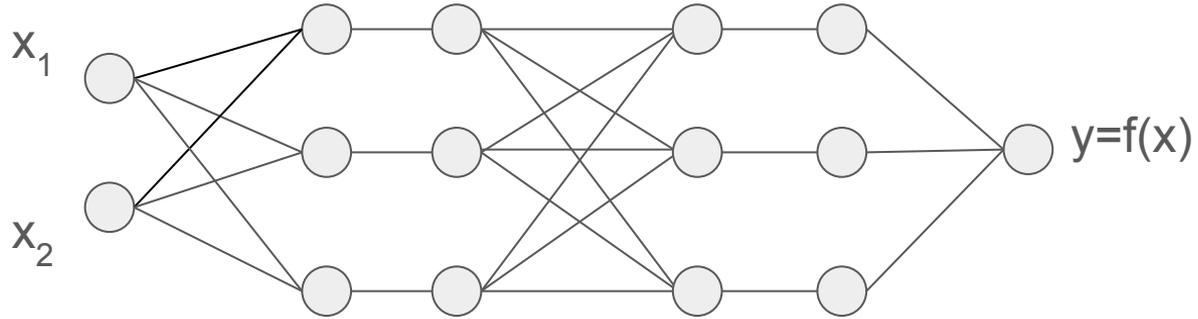
- Interval bound propagation (IBP)
- Linear (symbolic) bound propagation (CROWN)

Efficient methods are typically incomplete (solving a lower bound, as tight as possible)

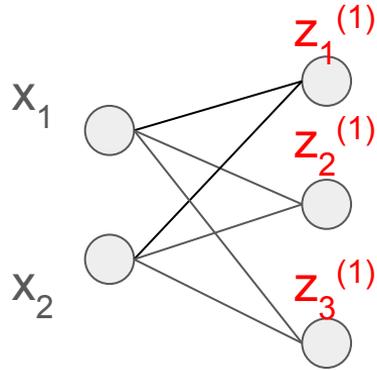


$$y^* = \min_{x \in \mathcal{S}} f(x)$$

Any faster ways to calculate the bounds on  $f(x)$ ?



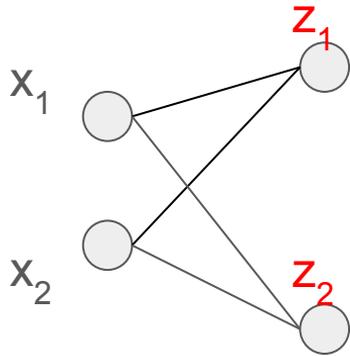
Let's look at one layer first



Given bounds on  $x$ , can we calculate the bounds on  $z$ ?

$$x_1 \in [-1, 2], x_2 \in [-2, 1]$$

## Let's look at one layer first



Given bounds on  $x$ , can we calculate the bounds on  $z$ ?

$$x_1 \in [-1, 2], x_2 \in [-2, 1]$$

As an illustration, suppose we have

$$z_1 = x_1 - x_2$$

$$z_2 = 2x_1 - x_2$$

Can you infer bounds on  $z$  given bounds on  $x$ ?

# Interval Bound Propagation (IBP)

$$x_1 \in [-1, 2], x_2 \in [-2, 1]$$

$$z_1 = x_1 - x_2$$

$$z_2 = 2x_1 - x_2$$

Lower bounds

$$\underline{z}_1 = -1 - 1 = -2$$

$$\underline{z}_2 = -1 \times 2 - 1 = -3$$

Upper bounds

$$\bar{z}_1 = 2 - (-2) = 4$$

$$\bar{z}_2 = 2 \times 2 - (-2) = 6$$

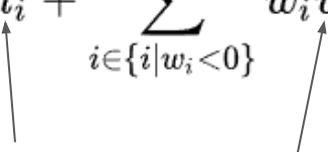
# Interval Bound Propagation (IBP)

$$x_1 \in [-1, 2], x_2 \in [-2, 1] \quad z_1 = x_1 - x_2 \quad z_2 = 2x_1 - x_2$$

$$\underline{z}_1 = -1 - 1 = -2 \quad \bar{z}_1 = 2 - (-2) = 4$$

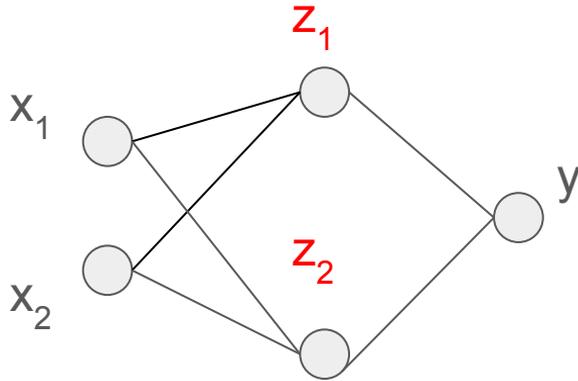
$$\underline{z}_2 = -1 \times 2 - 1 = -3 \quad \bar{z}_2 = 2 \times 2 - (-2) = 6$$

In general:

$$\sum_{i \in \{i | w_i \geq 0\}} w_i l_i + \sum_{i \in \{i | w_i < 0\}} w_i u_i \leq \sum_i w_i x_i \leq \sum_{i \in \{i | w_i \geq 0\}} w_i u_i + \sum_{i \in \{i | w_i < 0\}} w_i l_i$$


Elements lower and upper bounds of x

# Interval Bound Propagation: continue to the next layer



Let's say  $y = z_1 - z_2$

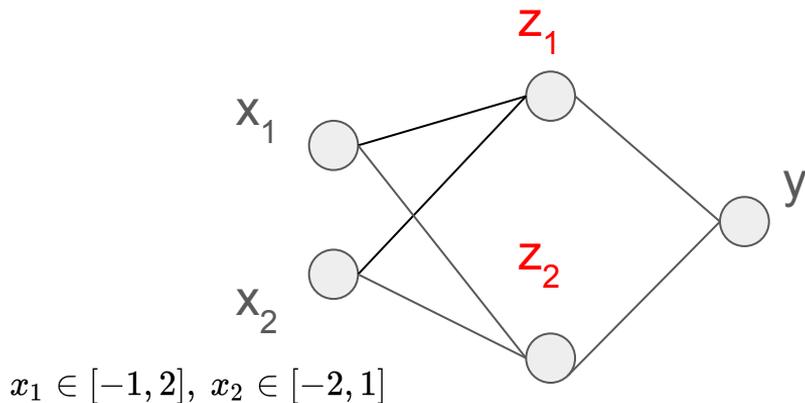
We also know that:

$$z_1 \in [-2, 4] \quad z_2 \in [-3, 6]$$

The what can we conclude about  $y$ ?

$$y \in [-8, 7]$$

# Interval Bound Propagation: limitations



Apply IBP we obtain  $y \in [-8, 7]$   
for this simple linear network.

However observe that

$$z_1 = x_1 - x_2$$

$$z_2 = 2x_1 - x_2$$

$$y = z_1 - z_2$$

$$y = x_1 - x_2 - (2x_1 - x_2) = -x_1$$

The actual bounds is  $[-2, 1]$ , much tighter than  $[-8, 7]$

## A Better Idea: Keep the correlations between $x$ and $z$

$$z_1 = x_1 - x_2$$

$$z_2 = 2x_1 - x_2$$

$$y = z_1 - z_2$$

$$y = x_1 - x_2 - (2x_1 - x_2) = -x_1$$

The actual bounds is  $[-2, 1]$ , much tighter than  $[-8, 7]$

It is important to keep the correlations between  $z$  and  $x$  to obtain this tighter result!

We treat  $z$  as a **symbolic function of  $x$** , rather than intervals

## A Better Idea: linear bound propagation

$$z_1 = x_1 - x_2$$

$$z_2 = 2x_1 - x_2$$

$$y = z_1 - z_2$$

$$y = x_1 - x_2 - (2x_1 - x_2) = -x_1$$

The actual bounds is  $[-2, 1]$ , much tighter than  $[-8, 7]$

It is important to keep the correlations between  $z$  and  $x$  to obtain this tighter result!

We treat  $z$  as a **linear function of  $x$** , rather than intervals

## A Better Idea: linear bound propagation

$$y = z_1 - z_2 \longrightarrow y = x_1 - x_2 - (2x_1 - x_2) = -x_1$$

Plug in

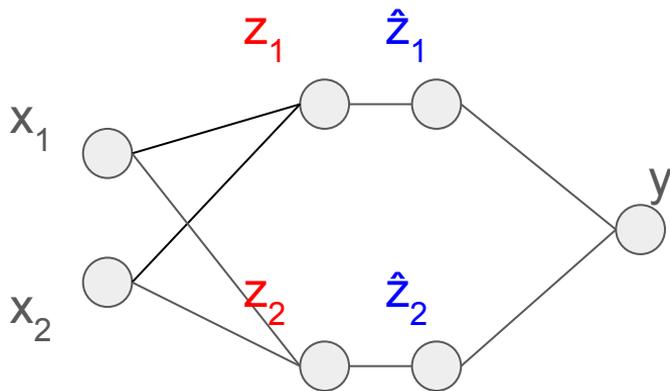
$$z_1 = x_1 - x_2$$

$$z_2 = 2x_1 - x_2$$

We treat  $z$  as a **linear function of  $x$** , rather than concrete intervals.

After we plug in linear functions ( $z$  w.r.t.  $x$ ), we still get a linear function ( $y$  w.r.t.  $x$ )

# Bound propagation: how about nonlinear functions?



Can we improve IBP using symbolic linear bounds?

Instead of  $y = z_1 - z_2$

Now we have  $y = \text{ReLU}(z_1) - \text{ReLU}(z_2)$

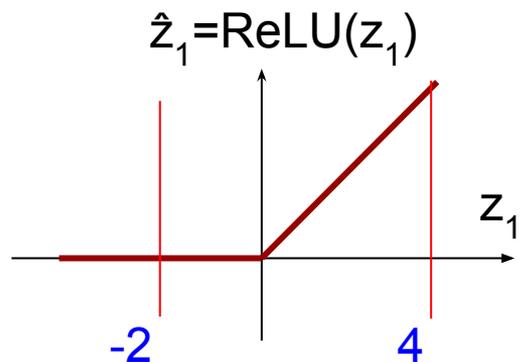
From IBP we already know that

$$z_1 \in [-2, 4], \quad z_2 \in [-3, 6],$$

$$\text{ReLU}(z_1) \in [0, 4], \quad \text{ReLU}(z_2) \in [0, 6]$$

$$y \in [-6, 4]$$

# Linear bound propagation for ReLU function (CROWN)



Instead of  $y = z_1 - z_2$

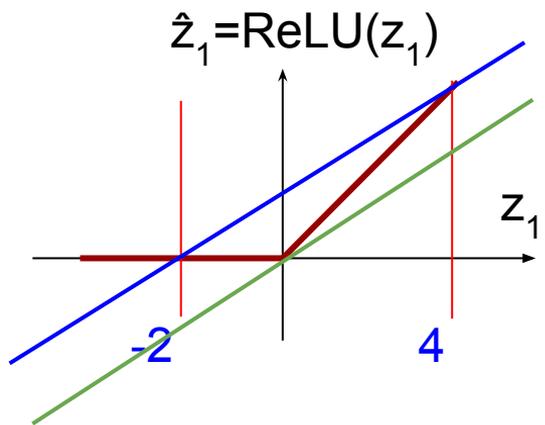
Now we have  $y = \text{ReLU}(z_1) - \text{ReLU}(z_2)$

We already know that

$$z_1 \in [-2, 4], \quad z_2 \in [-3, 6],$$

(Preactivation bounds)

# Linear bound propagation for ReLU function (CROWN)

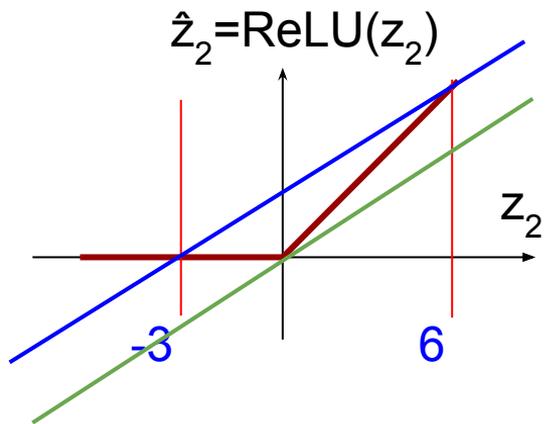


Linear **upper bound** (same as the one of triangle relaxation in LP)

Linear **lower bound** (actually not unique)

$$\boxed{\frac{2}{3}z_1} \leq \text{ReLU}(z_1) \leq \boxed{\frac{2}{3}z_1 + \frac{4}{3}}$$

# Linear bound propagation for ReLU function (CROWN)



$\text{ReLU}(z_2)$  can be bounded using linear functions similarly.

Now let's consider  $y = \text{ReLU}(z_1) - \text{ReLU}(z_2)$ . How to bound it using linear functions of  $z_1$  and  $z_2$ ?

$$\frac{2}{3}z_1 \leq \text{ReLU}(z_1) \leq \frac{2}{3}z_1 + \frac{4}{3}$$

$$\frac{2}{3}z_2 \leq \text{ReLU}(z_2) \leq \frac{2}{3}z_2 + 2$$

# Linear bound propagation for ReLU function (CROWN)

$$\boxed{\frac{2}{3}z_1} \leq \text{ReLU}(z_1) \leq \frac{2}{3}z_1 + \frac{4}{3}$$

$$\frac{2}{3}z_2 \leq \text{ReLU}(z_2) \leq \boxed{\frac{2}{3}z_2 + 2}$$

Negative coefficient, take  
upper bound

$$\boxed{\frac{2}{3}z_1} - \boxed{\left(\frac{2}{3}z_2 + 2\right)} \leq$$

$$y = \text{ReLU}(z_1) - \text{ReLU}(z_2)$$

positive coefficient, take  
lower bound

$$\leq \left(\frac{2}{3}z_1 + \frac{4}{3}\right) - \frac{2}{3}z_2$$

## Linear bound propagation for ReLU function (CROWN)

$$\frac{2}{3}z_1 - \left(\frac{2}{3}z_2 + 2\right) \leq y \leq \left(\frac{2}{3}z_1 + \frac{4}{3}\right) - \frac{2}{3}z_2$$

Now we have linear inequalities for  $y$  w.r.t.  $z$ !

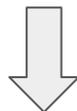
Next step we can simply plug in, as in the linear ( $y=z_1-z_2$ ) case.

# Linear bound propagation for ReLU function (CROWN)

$$\frac{2}{3}z_1 - \left(\frac{2}{3}z_2 + 2\right) \leq y \leq \left(\frac{2}{3}z_1 + \frac{4}{3}\right) - \frac{2}{3}z_2$$

$$z_1 = x_1 - x_2$$

$$z_2 = 2x_1 - x_2$$



Plug in

$$-\frac{2}{3}x_1 - 2 \leq y \leq -\frac{2}{3}x_1 + \frac{4}{3}$$

# Linear bound propagation for ReLU function (CROWN)

We now have symbolic linear bounds for  $y$  w.r.t.  $x$

$$\boxed{-\frac{2}{3}x_1 - 2} \leq y \leq \boxed{-\frac{2}{3}x_1 + \frac{4}{3}}$$

$$x_1 \in [-1, 2], x_2 \in [-2, 1]$$

Take lower bound given  $x$

Take upper bound given  $x$

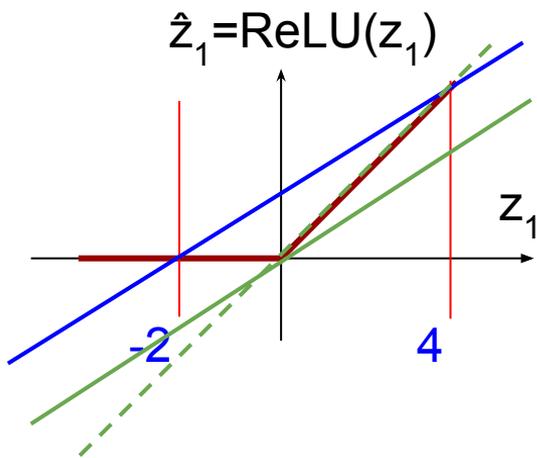
$$y \in \left[-\frac{10}{3}, 2\right]$$

Concrete interval bounds

A lot more tighter than IBP bounds  $y \in [-6, 4]$

# Can we do even better?

Let's recall that when we linearly bound the ReLU function, there are some flexibilities



Linear **upper bound** (same as the one of triangle relaxation in LP)

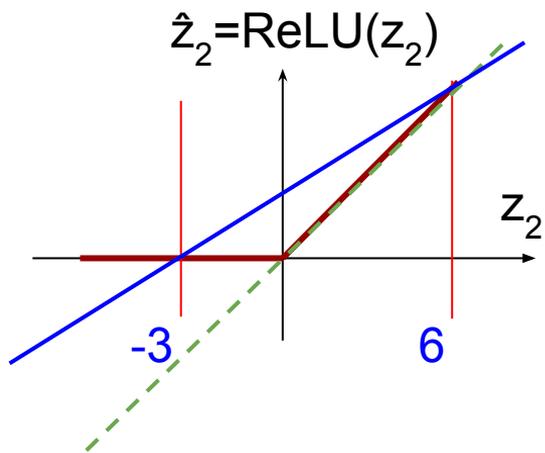
Linear **lower bound** (actually **not unique**)

$$\frac{2}{3}z_1 \leq \text{ReLU}(z_1) \leq \frac{2}{3}z_1 + \frac{4}{3}$$

Also valid:  $z_1 \leq \text{ReLU}(z_1) \leq \frac{2}{3}z_1 + \frac{4}{3}$

# Choosing different linear bounds ( $\alpha$ -CROWN)

Now what are the linear bounds of  
 $y = \text{ReLU}(z_1) - \text{ReLU}(z_2)$ ?



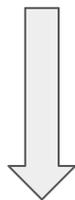
$$z_1 - \left(\frac{2}{3}z_2 + 2\right) \leq y \leq \left(\frac{2}{3}z_1 + \frac{4}{3}\right) - z_2$$

$$z_1 \leq \text{ReLU}(z_1) \leq \frac{2}{3}z_1 + \frac{4}{3}$$

$$z_2 \leq \text{ReLU}(z_2) \leq \frac{2}{3}z_2 + 2$$

## Choosing different linear bounds ( $\alpha$ -CROWN)

$$z_1 - \left(\frac{2}{3}z_2 + 2\right) \leq y \leq \left(\frac{2}{3}z_1 + \frac{4}{3}\right) - z_2$$

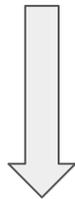


$$z_1 = x_1 - x_2$$

$$z_2 = 2x_1 - x_2$$

Plug in

$$-\frac{1}{3}x_1 - \frac{1}{3}x_2 - 2 \leq y \leq -\frac{4}{3}x_1 + \frac{1}{3}x_2 + \frac{4}{3}$$

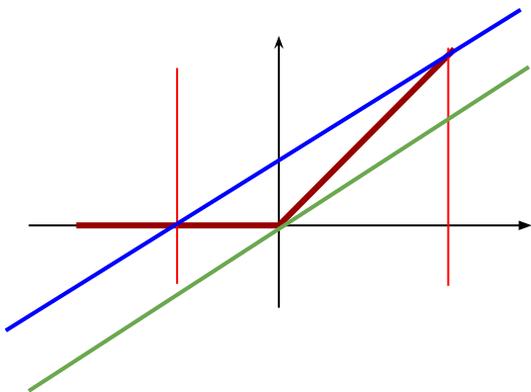


$$x_1 \in [-1, 2], x_2 \in [-2, 1]$$

Concretize

$$y \in [-3, 3]$$

# Linear lower bounds for ReLU function matters!

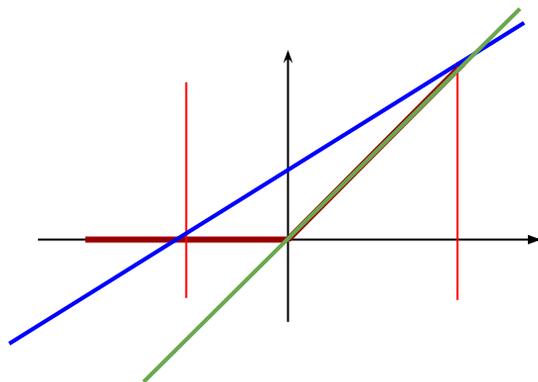


$$\frac{2}{3}z_1 \leq \text{ReLU}(z_1) \leq \frac{2}{3}z_1 + \frac{4}{3}$$

$$\frac{2}{3}z_2 \leq \text{ReLU}(z_2) \leq \frac{2}{3}z_2 + 2$$

↓

$$y \in \left[-\frac{10}{3}, 2\right]$$



$$z_1 \leq \text{ReLU}(z_1) \leq \frac{2}{3}z_1 + \frac{4}{3}$$

$$z_2 \leq \text{ReLU}(z_2) \leq \frac{2}{3}z_2 + 2$$

↓

$$y \in [-3, 3]$$

Which one is correct?

# Linear lower bounds for ReLU function matters!

Both results are correct! But we want the bounds to be as tight as possible! So best result is  $\mathbf{y} \in [-3, 2]$

In general, the slope of the linear lower bound for every ReLU neuron can be optimized to find the best result.

# Linear lower bounds for ReLU function matters!

In general, the slope of the linear lower bound for every ReLU neuron can be optimized to find the best result.

$$\alpha_1 z_2 \leq \text{ReLU}(z_2) \leq \frac{2}{3} z_2 + \frac{4}{3}$$

$$\alpha_2 z_2 \leq \text{ReLU}(z_2) \leq \frac{2}{3} z_2 + 2$$

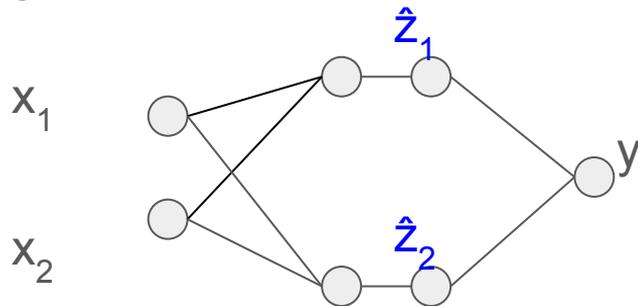
For optimal lower bound of  $y$ , set  $\alpha_1=1$ ,  $\alpha_2=1$

For optimal upper bound of  $y$ , set  $\alpha_1=2/3$ ,  $\alpha_2=2/3$

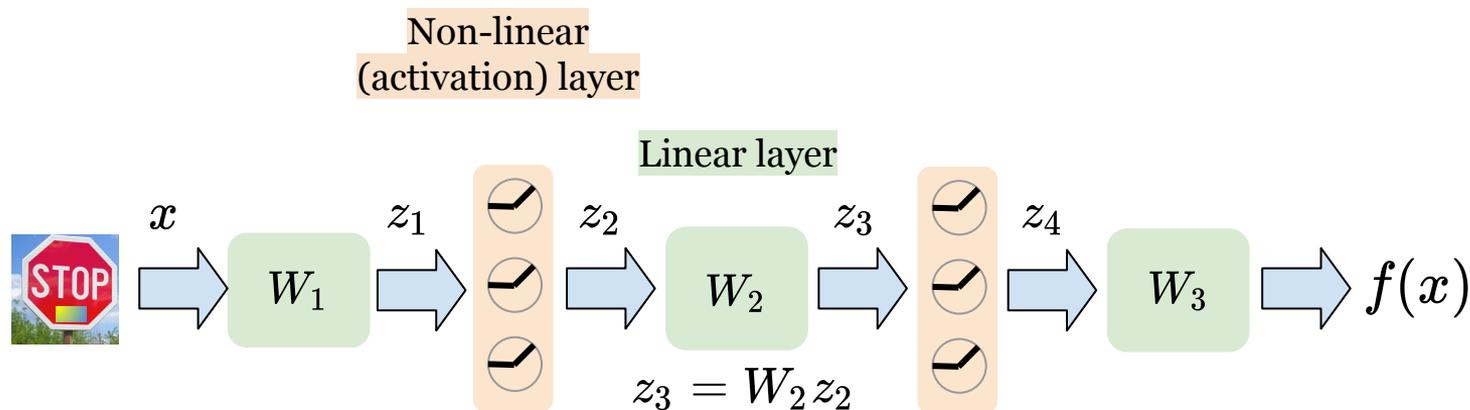
(note that the optimal  $\alpha_1$  and  $\alpha_2$  do not equal in general)

# Linear bound propagation method (CROWN)

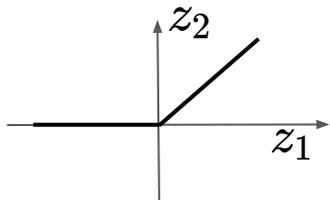
1. Obtain all pre-activation bounds (can be done via CROWN recursively)
2. Start from the output layer, form the initial linear (in)equality  $y = y$
3. Recursively propagate linear inequality  $y \leq a^T z + b$  through each layer:
  - a. For a linear layer,  $z = Wz'$ , directly plug in  $a^T z + b$  to get a linear bound of  $z'$
  - b. For a non-linear layer (e.g.,  $z = \text{ReLU}(z')$ ), we first form the linear inequalities to bound the nonlinear layer itself. Then multiply either the lower or upper bound based on the sign of element in  $a$
4. When the linear inequality propagates to the input layer, we can concretize the linear bound using bounds on input layer.



# Illustration: Linear bound propagation process



$$z_2 = \text{ReLU}(z_1)$$

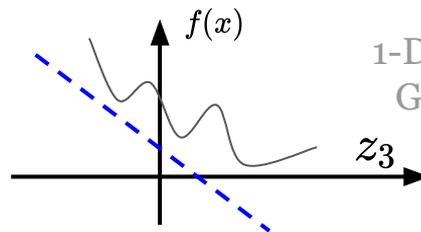


Steps:

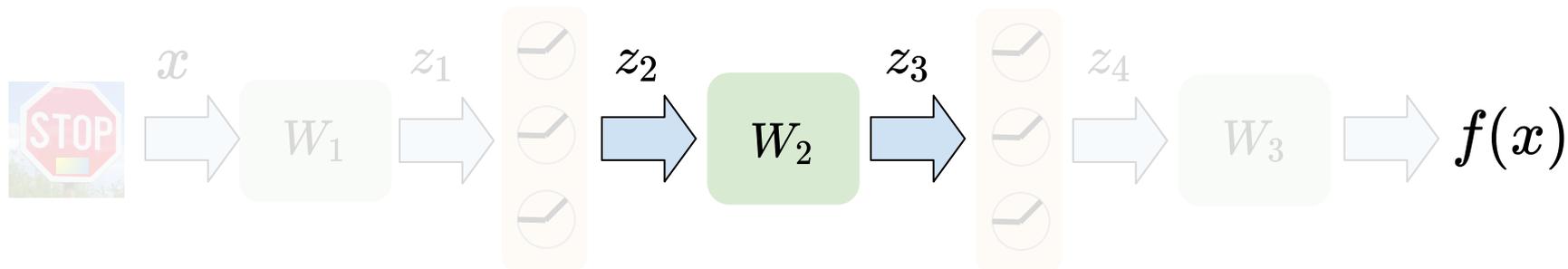
- Propagate bounds through linear layers
- Propagate bounds through non-linear layers

# Illustration: Linear bound propagation process

A linear lower bound for an intermediate layer

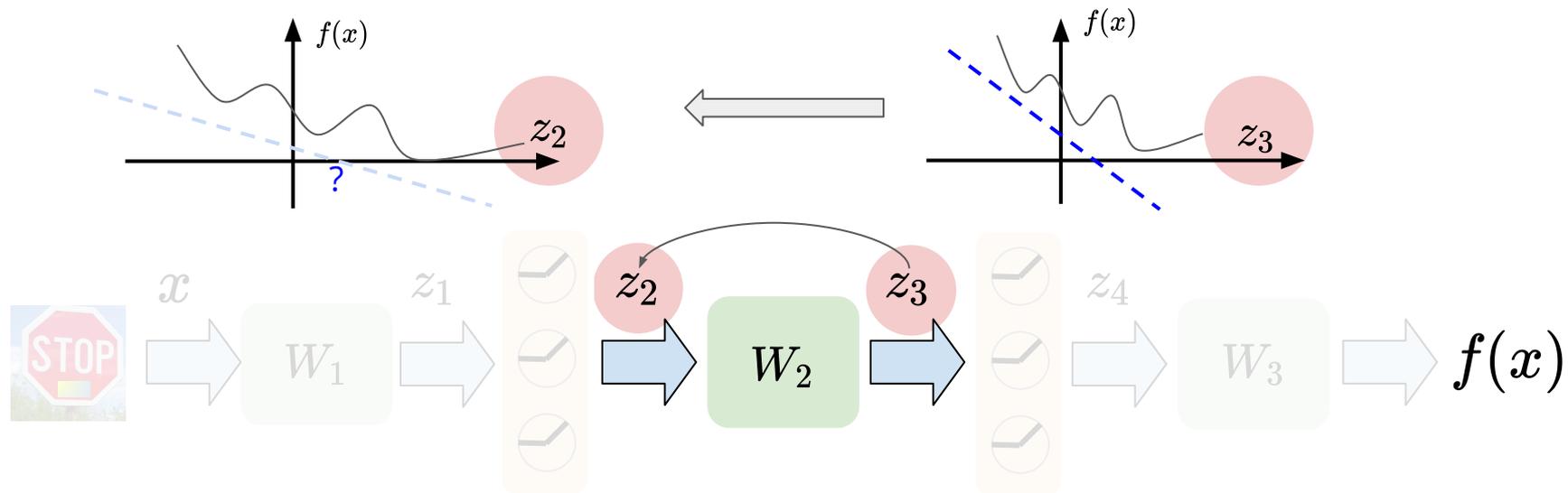


1-D case for illustration.  
Generally it's a linear hyperplane



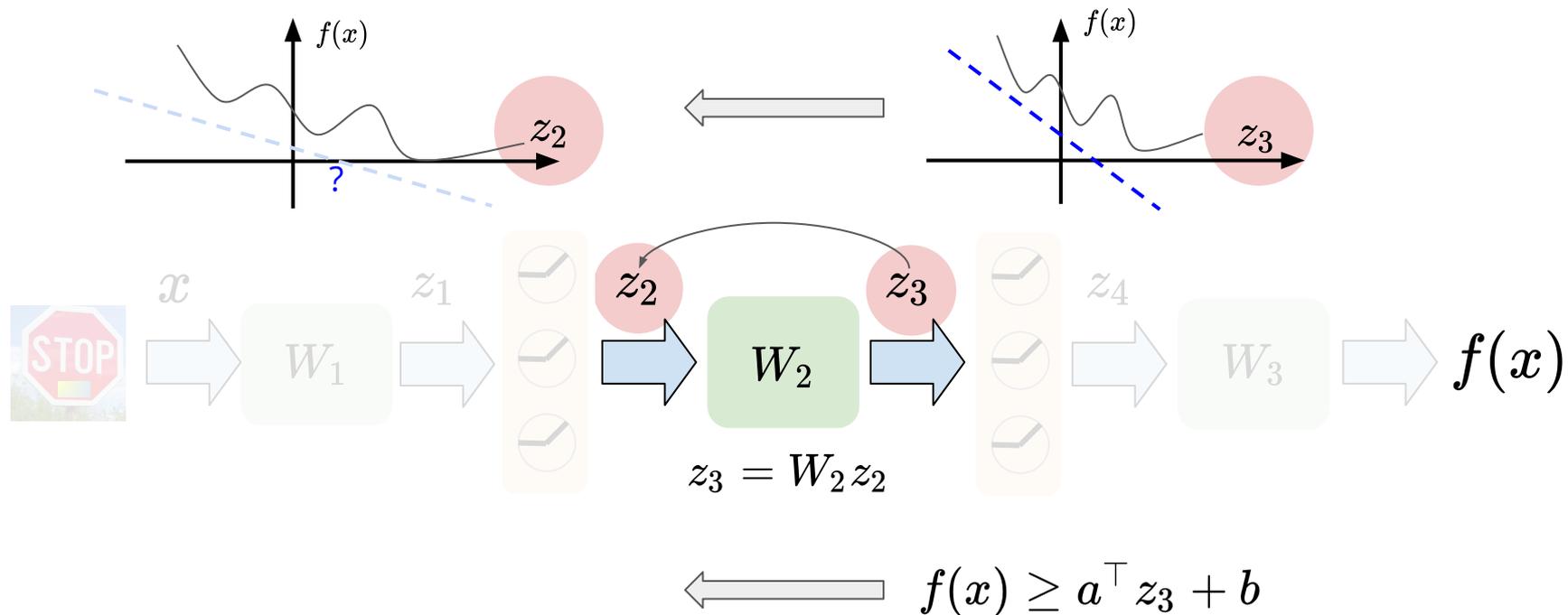
$$f(x) \geq a^\top z_3 + b$$

# Illustration: Linear bound propagation process

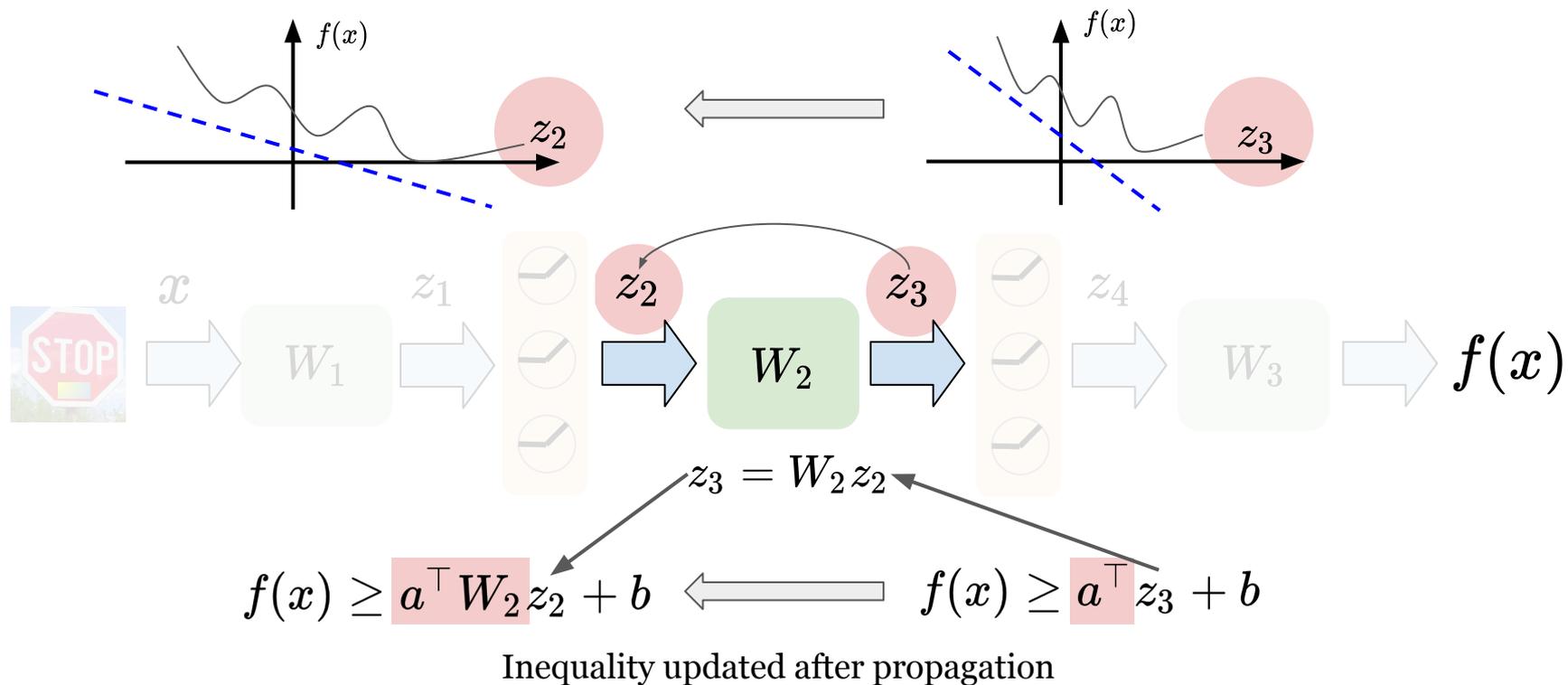


Propagate it to one layer before,  
while keeping the lower bound valid

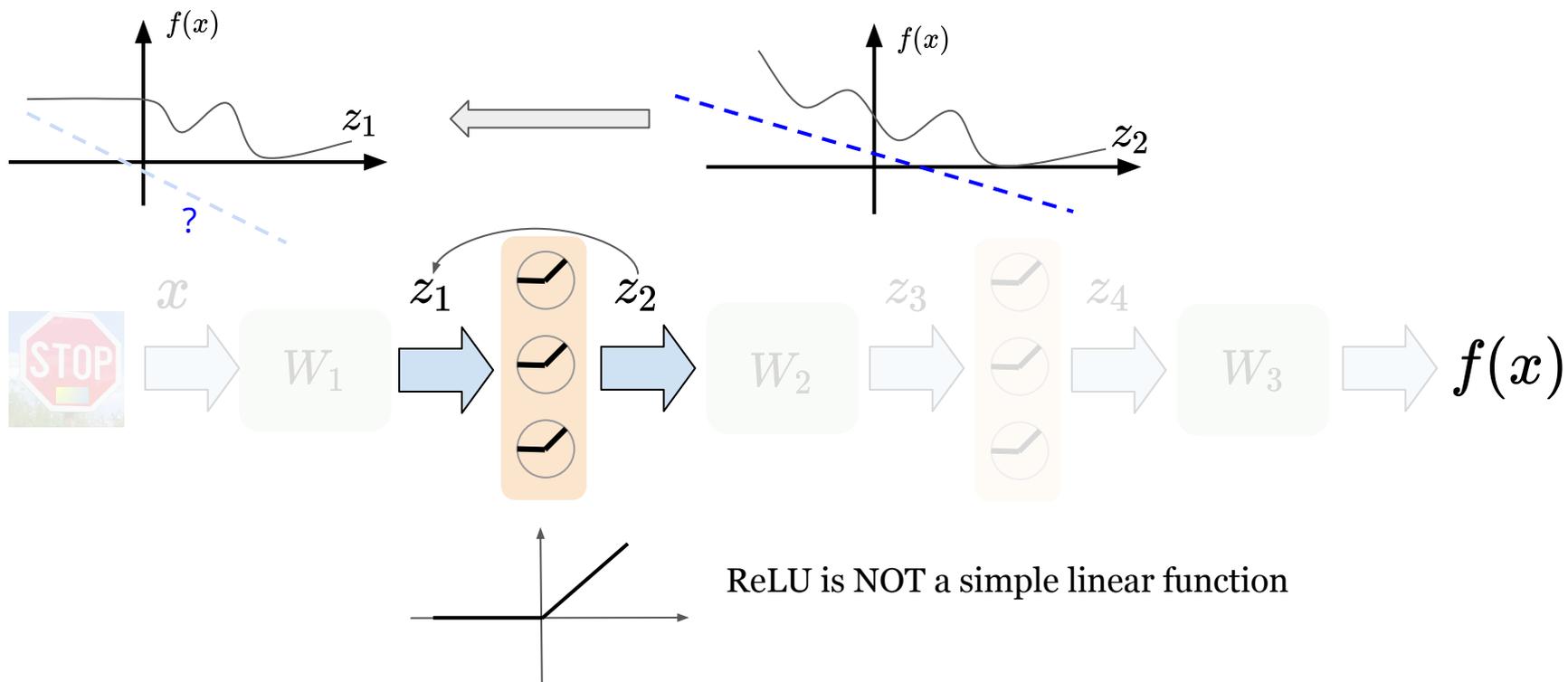
# Illustration: Linear bound propagation process



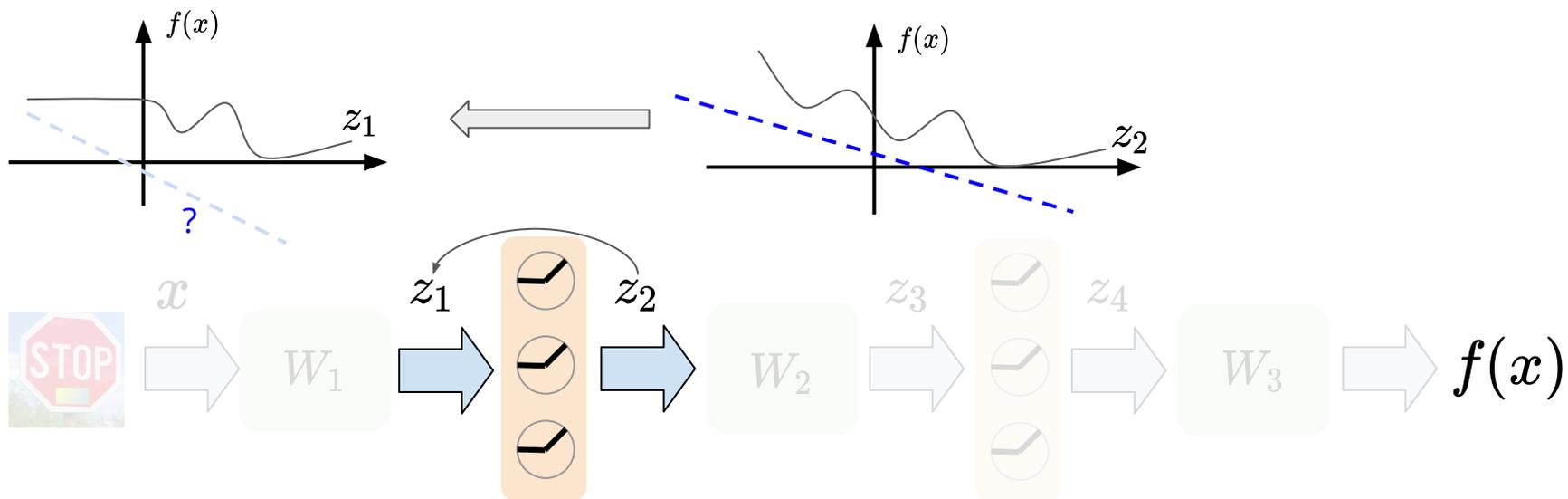
# Illustration: Linear bound propagation process



# Illustration: Linear bound propagation process

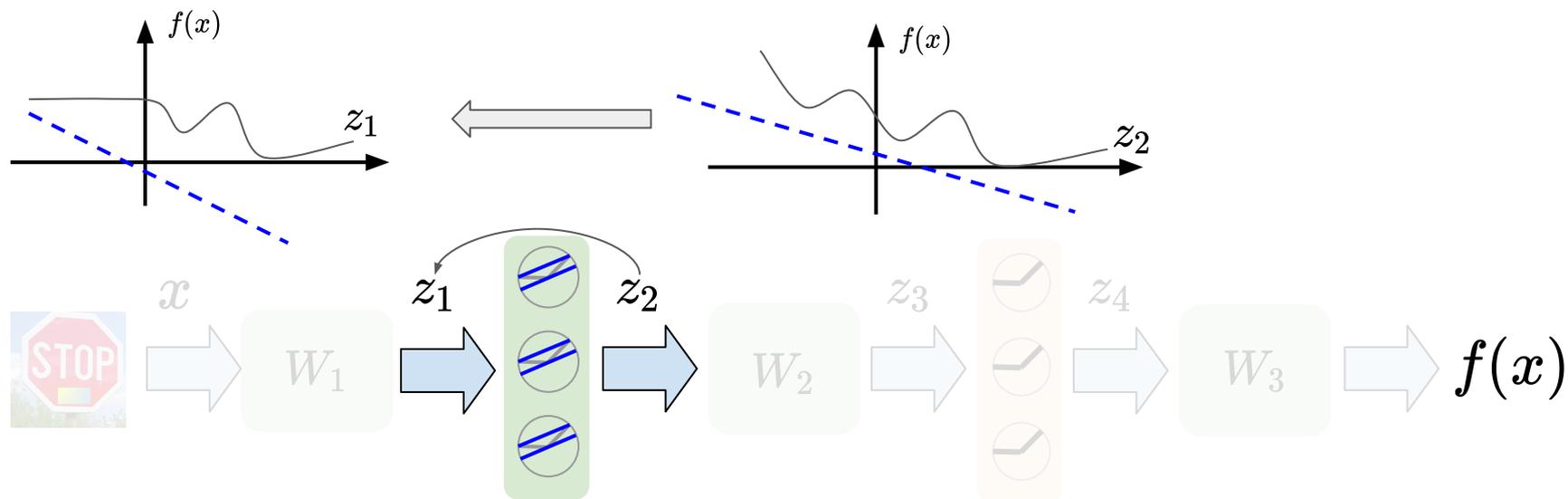


# Illustration: Linear bound propagation process



$$\longleftarrow f(x) \geq a^\top W_2 z_2 + b \quad \forall x \in \mathcal{S}$$

# Illustration: Linear bound propagation process

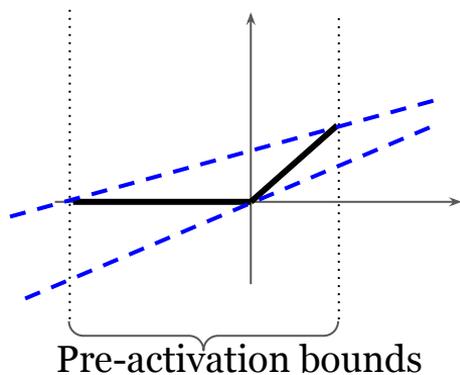


**Theorem** (informal): we can efficiently find  $D$ ,  $b'$  such that:

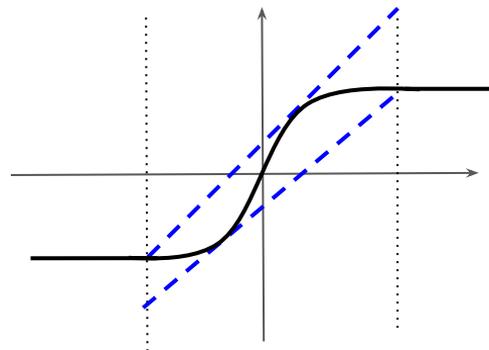
$$f(x) \geq a^\top W_2 D z_1 + b' \longleftarrow f(x) \geq a^\top W_2 z_2 + b \quad \forall x \in \mathcal{S}$$

# Illustration: Linear bound propagation process

**Proof sketch:** conservatively use **linear bounds** to **replace a non-linear function**.



(can be pre-computed using CROWN)

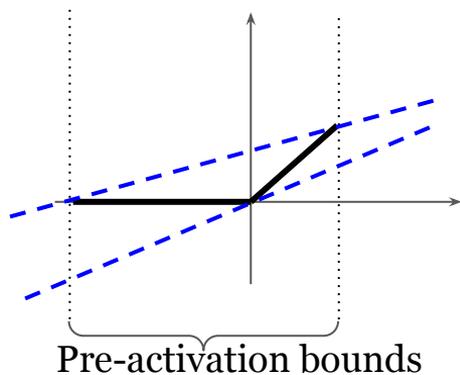


**Theorem** (informal): we can efficiently find  $D$ ,  $b'$  such that:

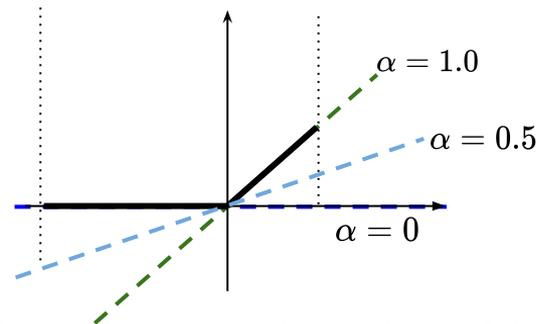
$$f(x) \geq a^\top W_2 D z_1 + b' \longleftarrow f(x) \geq a^\top W_2 z_2 + b \quad \forall x \in \mathcal{S}$$

# Illustration: Linear bound propagation process

Proof sketch: conservatively use linear bounds to replace a non-linear function.



(can be pre-computed using CROWN)



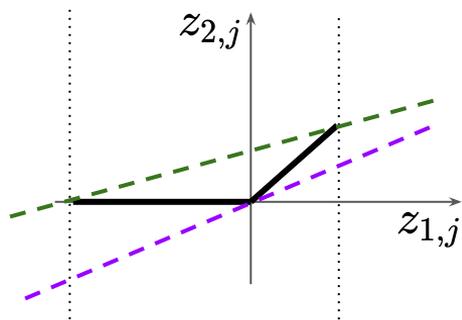
ReLU's lower bound can be optimized ( $\alpha$ -CROWN)

**Theorem** (informal): we can efficiently find  $D$ ,  $b'$  such that:

$$f(x) \geq a^\top W_2 D z_1 + b' \longleftarrow f(x) \geq a^\top W_2 z_2 + b \quad \forall x \in \mathcal{S}$$

# Illustration: Linear bound propagation process

Proof sketch: conservatively use **linear bounds** to **replace a non-linear function**.



$$f(x) \geq a^\top W_2 z_2 + b$$

$$f(x) \geq \sum_j [(a^\top W_2)_j \cdot z_{2,j}] + b$$

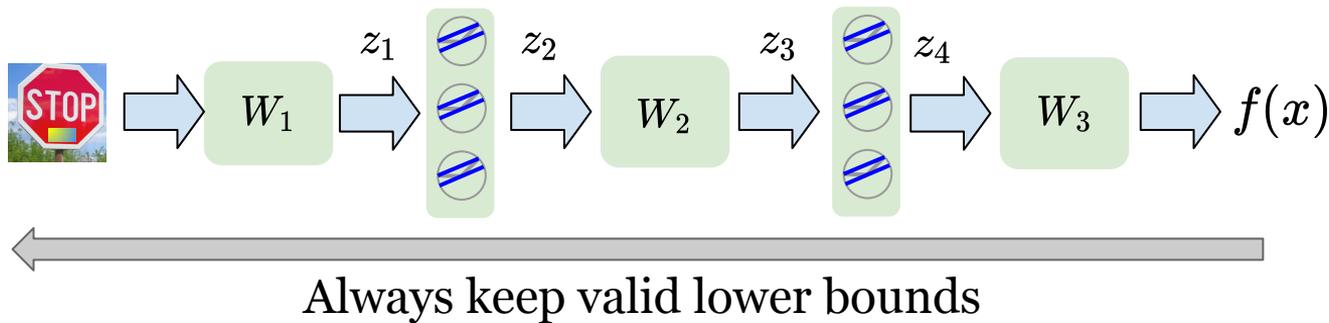
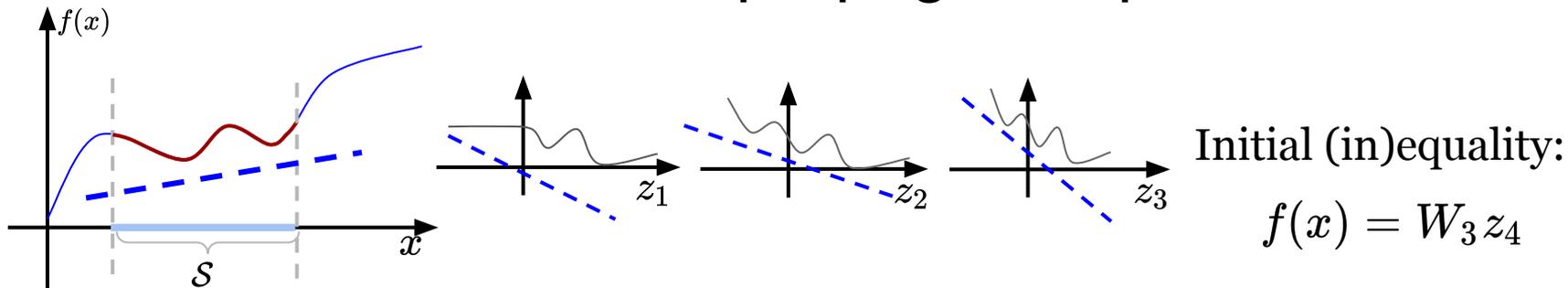
$(a^\top W_2)_j \geq 0$  Choose **lower bound**

$(a^\top W_2)_j < 0$  Choose **upper bound**

**Theorem** (informal): we can efficiently find  $D$ ,  $b'$  such that:

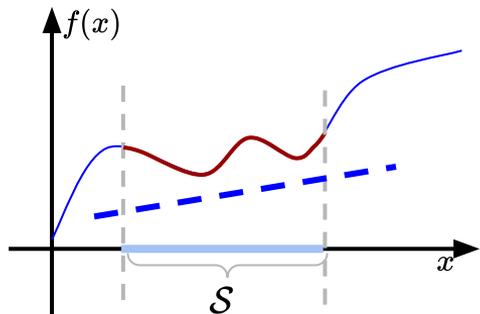
$$f(x) \geq a^\top W_2 D z_1 + b' \longleftarrow f(x) \geq a^\top W_2 z_2 + b \quad \forall x \in \mathcal{S}$$

# Illustration: Linear bound propagation process

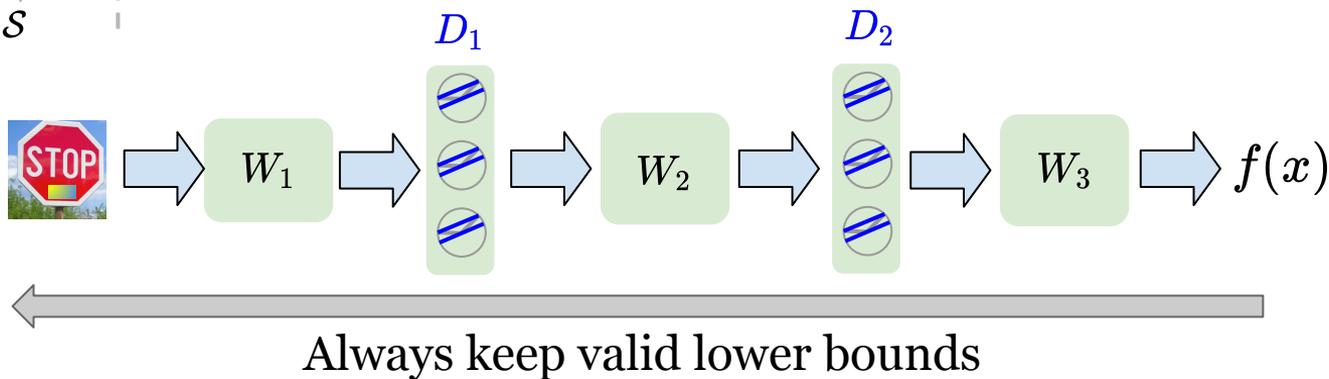


**CROWN main theorem** (simplified):  $f(x) \geq a_{\text{CROWN}}^\top x + b_{\text{CROWN}} \quad \forall x \in \mathcal{S}$

# Illustration: Linear bound propagation process



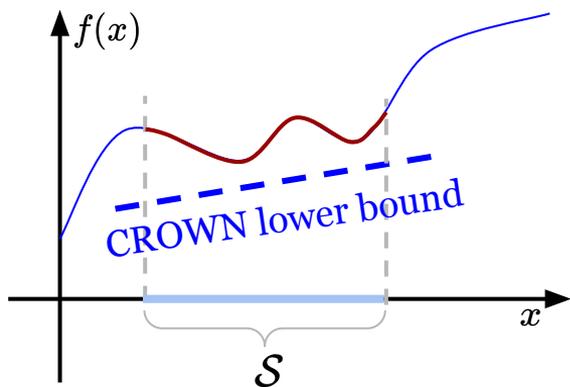
Bounds propagated through  
simple matrix multiplications!  
Fast and GPU-friendly



**CROWN main theorem** (simplified):  $f(x) \geq a_{\text{CROWN}}^\top x + b_{\text{CROWN}} \quad \forall x \in \mathcal{S}$

$$a_{\text{CROWN}} = W_3 D_2 W_2 D_1 W_1$$

# Use Linear Bounds to Prove Robustness



Prove:  $\forall x \in \mathcal{S}, f(x) > 0$



$x_1 \in \mathcal{S}$



$x_2 \in \mathcal{S}$



$x_3 \in \mathcal{S}$

...

Lower bound  $> 0 \implies f(x) > 0 \implies$  verified (always a stop sign)

# auto\_LiRPA: Verification Library for General Computation Graphs

 Colab Demo:

<http://PaperCode.cc/AutoLiRPA-Demo>



The auto\_LiRPA library on GitHub:

<http://PaperCode.cc/AutoLiRPA>

# MILP/LP vs Bound Propagation

Bound propagation:

- Scalable and fast propagation
- GPU friendly
- Incomplete verification (will be extended in the next lecture)
- Bounds are looser compared to LP; much looser compared to MILP

MILP/LP:

- Tighter solution
- Does not scale (MILP ~10k neurons, LP ~100k neurons)
- Much slower; cannot utilize GPU